

C G T A C G T A
A C G T A C G T

Whole Genome Alignment

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@aphillippy 

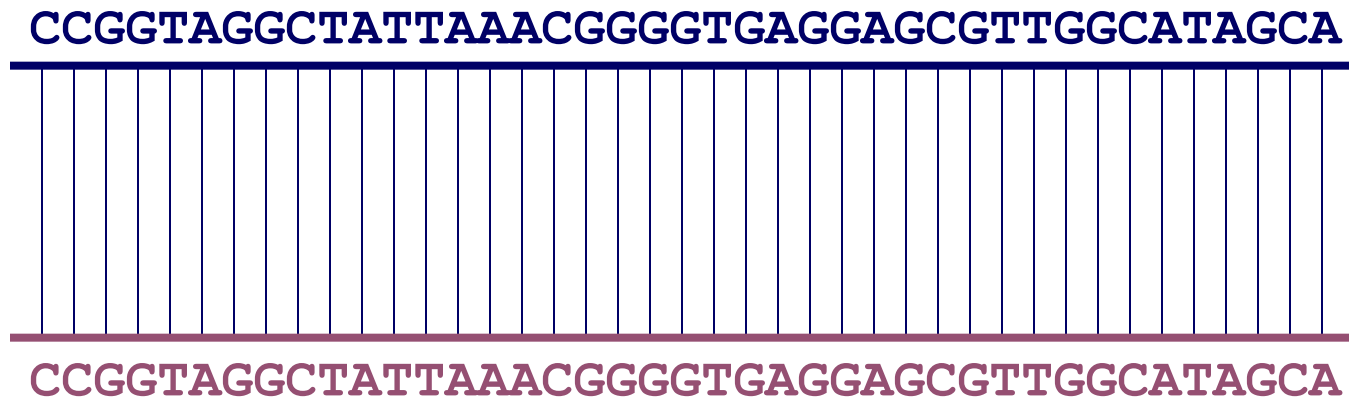


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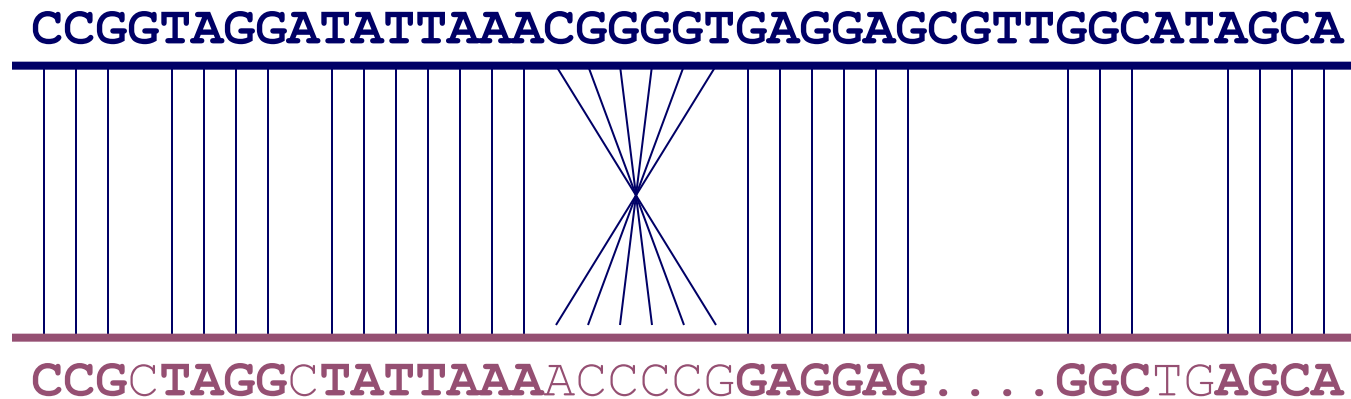
Goal of whole genome alignment

- For genomes A and B , find a mapping from each position in A to its corresponding position in B



Not so fast...

- Genome *A* may have insertions, deletions, translocations, inversions, or duplications with respect to *B*

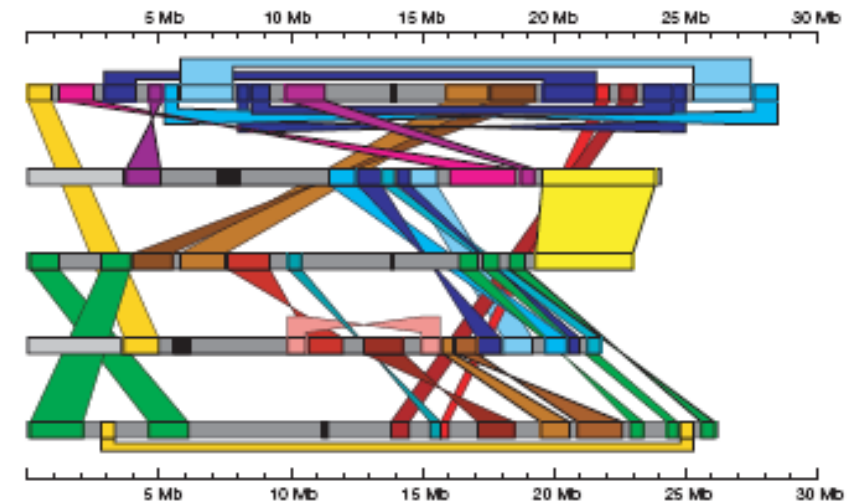


Global vs. local alignments

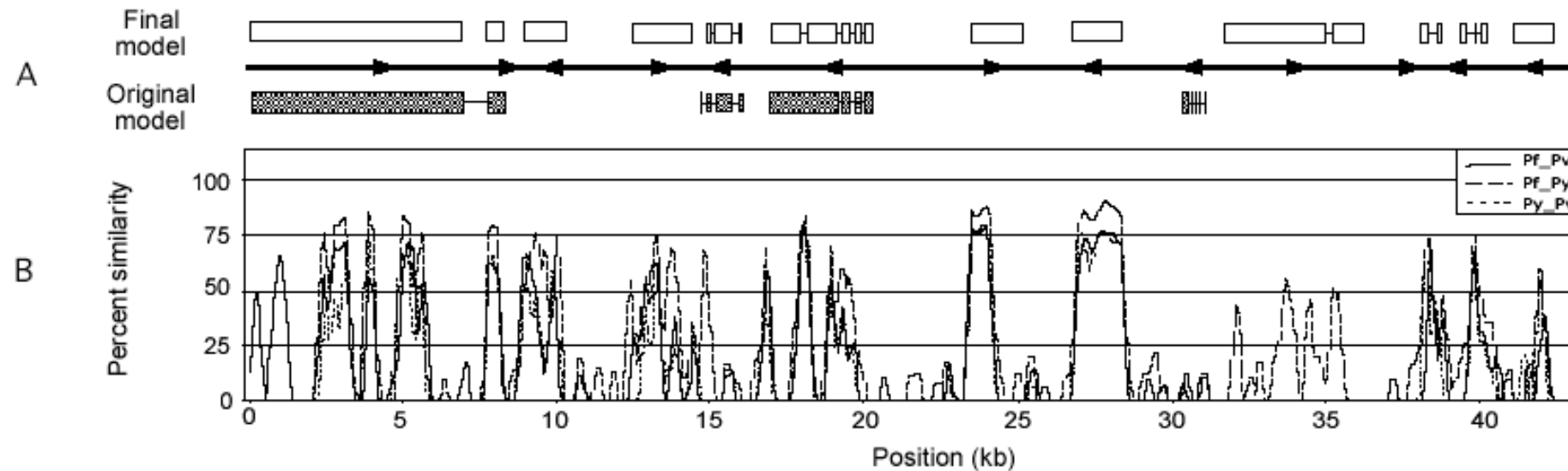
- Global pairwise alignment

...AAGCTTGGCTTAGCTGCTAGGGTAGGCTTGGG...
...AAGCTGGGCTTAGTTGCTAG..TAGGCTTTGG...
 ^ ^ ^^ ^

- Whole genome alignment
 - A “collection” of local alignments



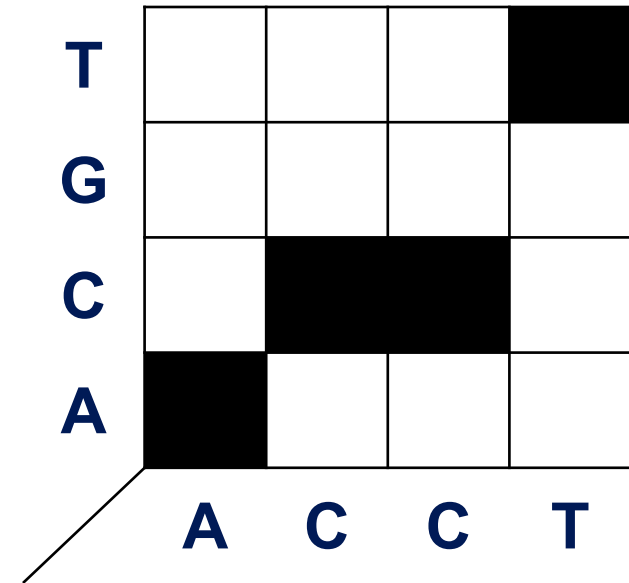
Global alignment visualization



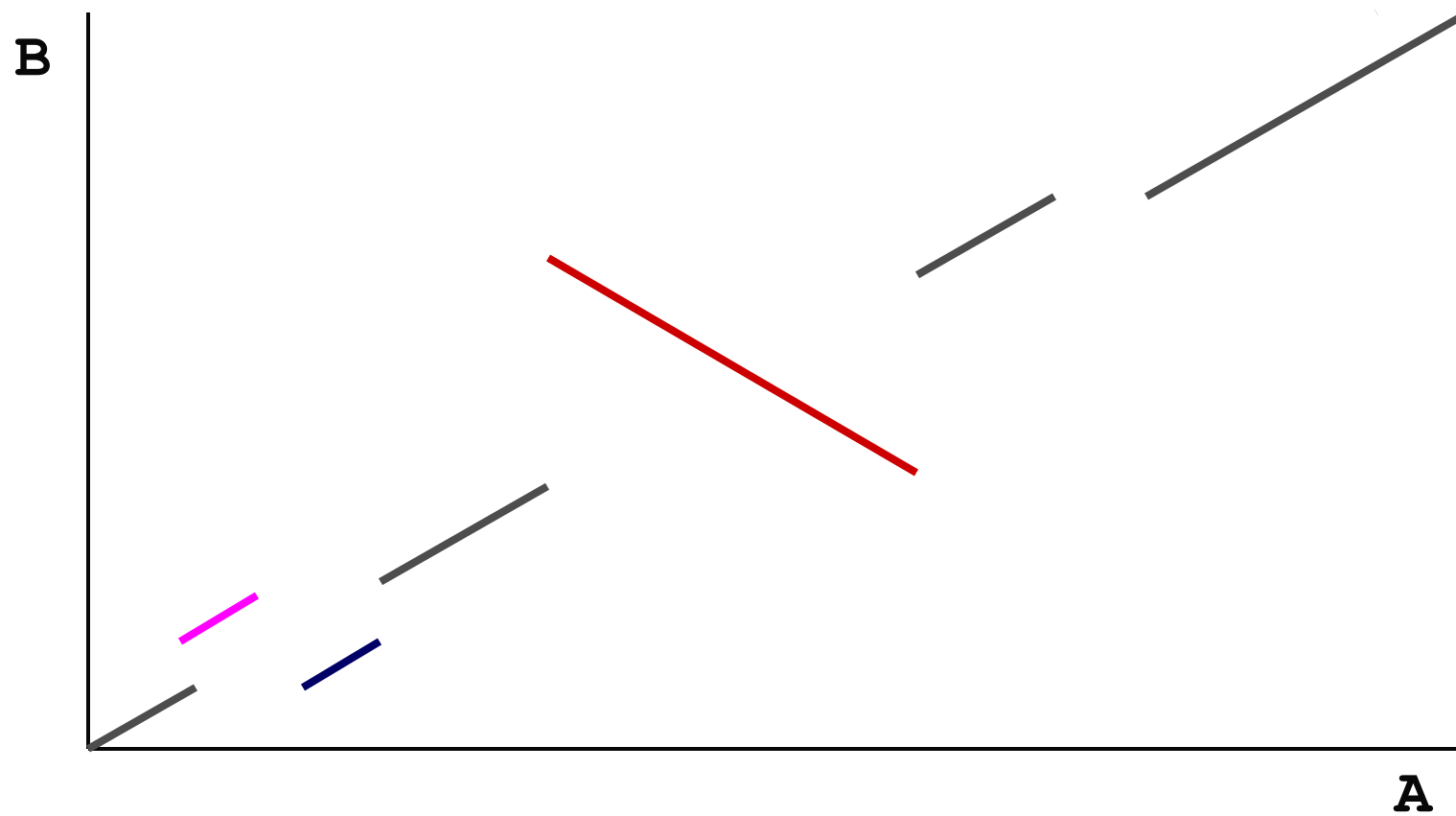
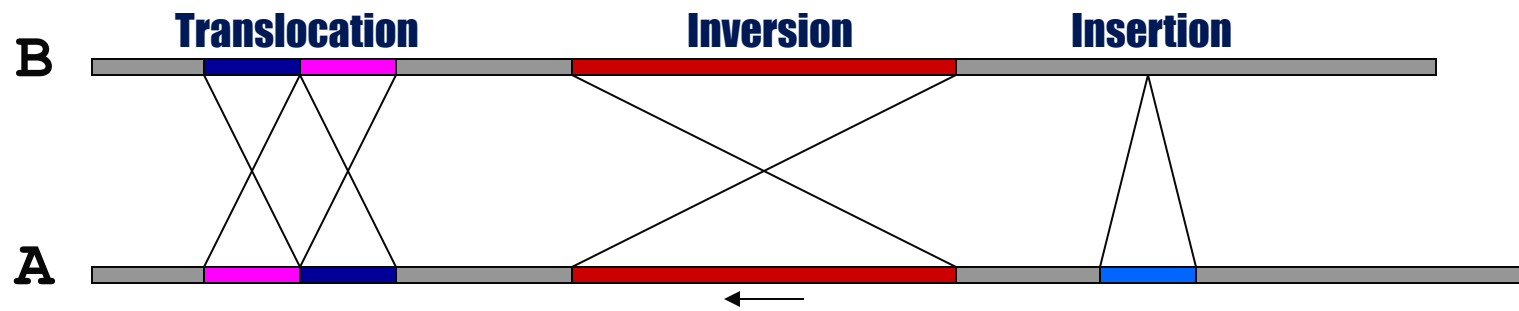
Whole genome alignment visualization

- How can we visualize *whole* genome alignments?

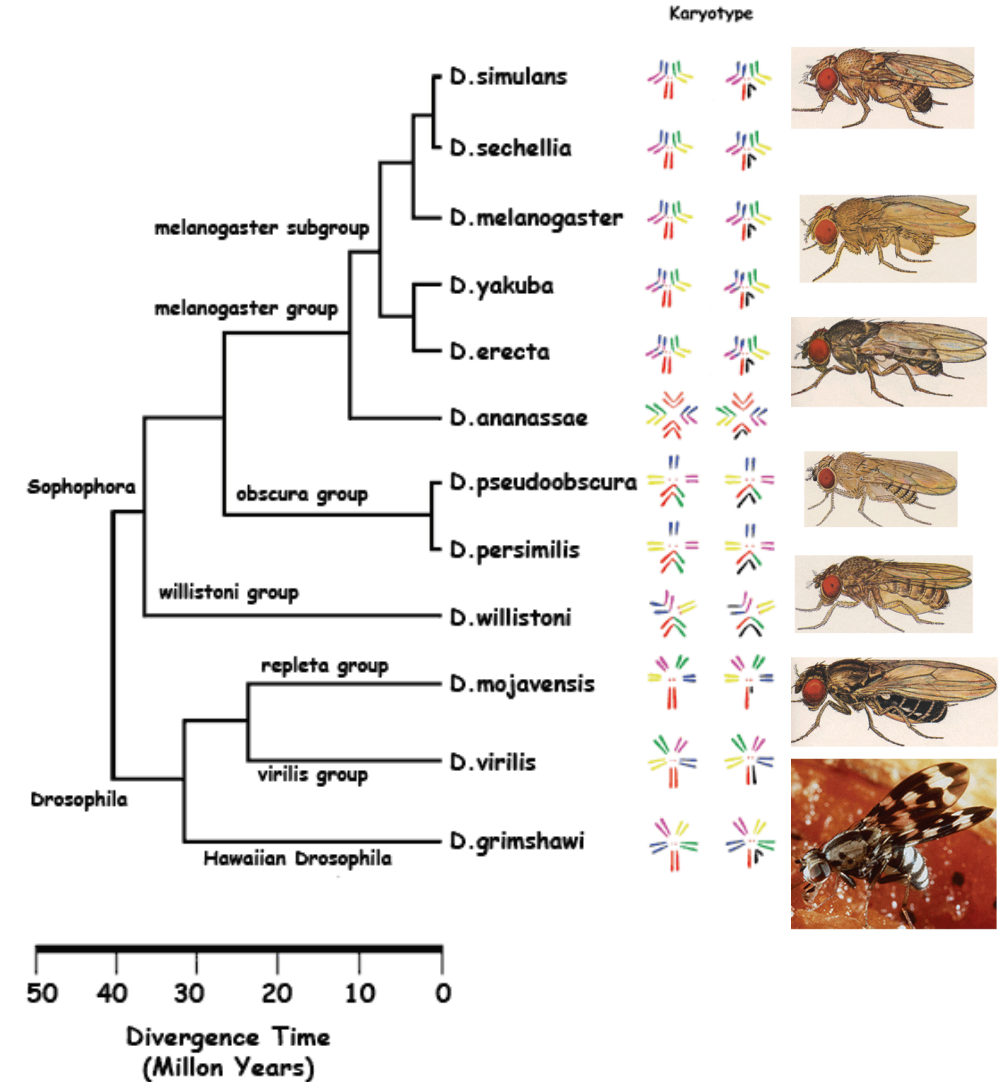
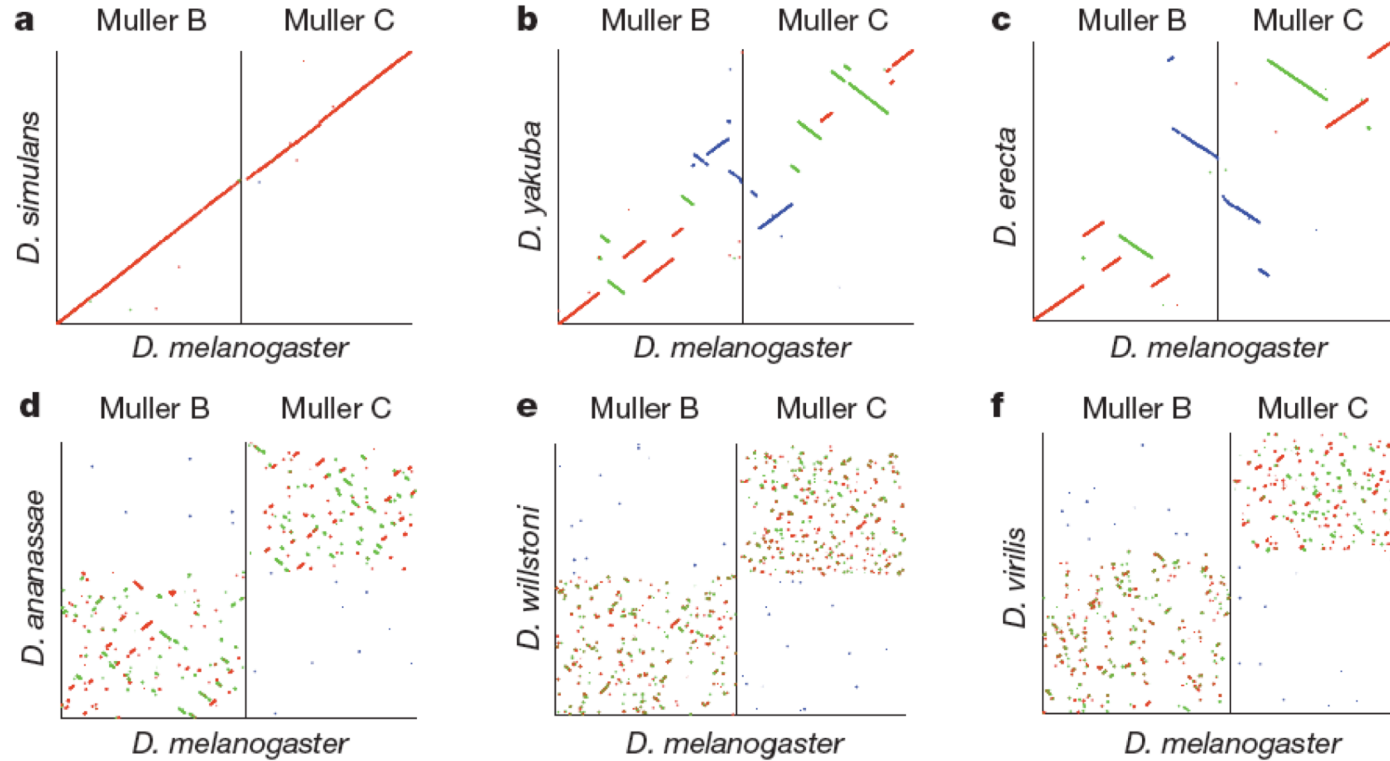
- With an alignment dot plot
 - $N \times M$ matrix
 - Let i = position in genome A
 - Let j = position in genome B
 - Fill cell (i,j) if A_i shows similarity to B_j



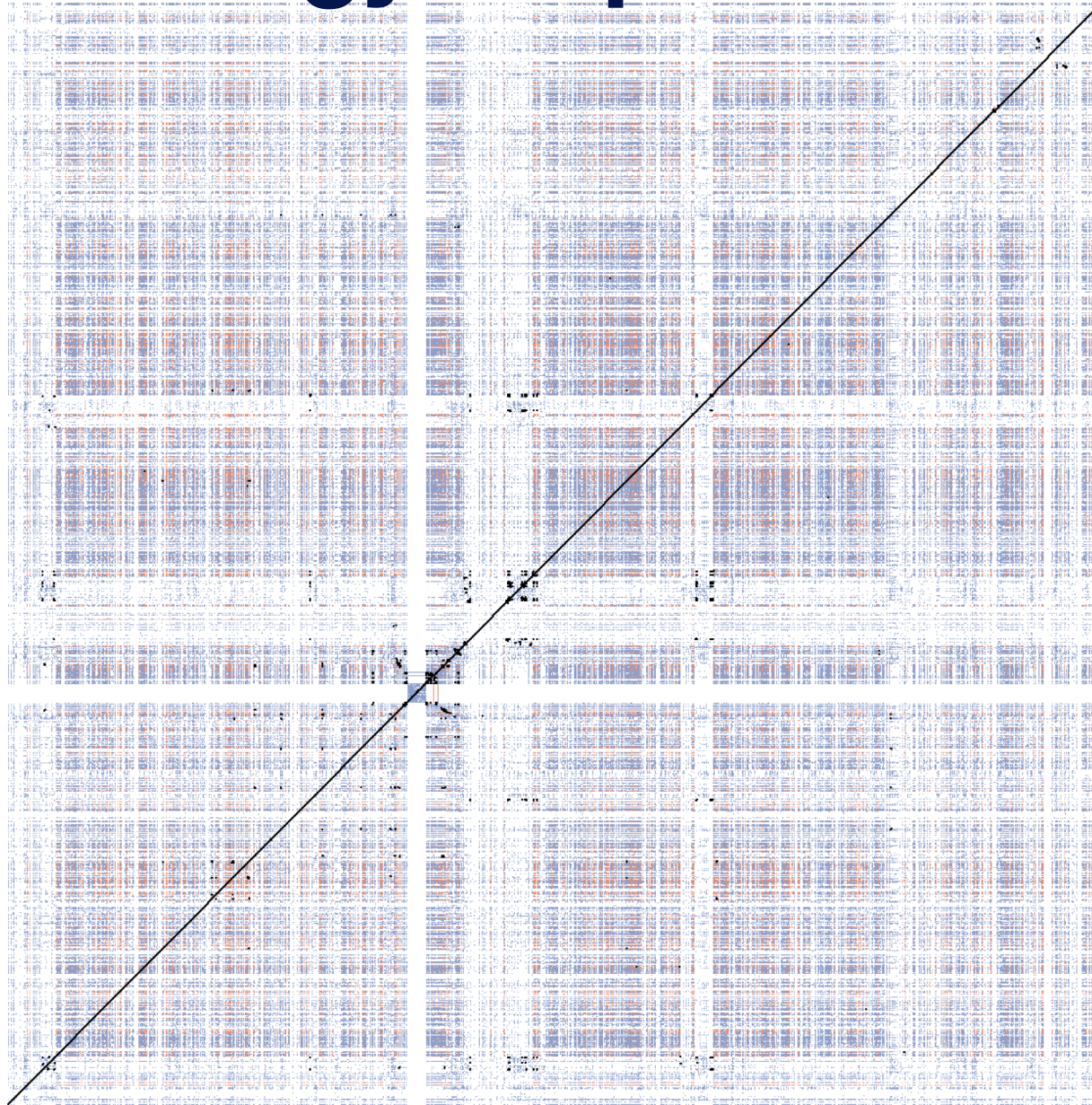
- A perfect alignment between A and B would completely fill the positive diagonal



Drosophila genome evolution

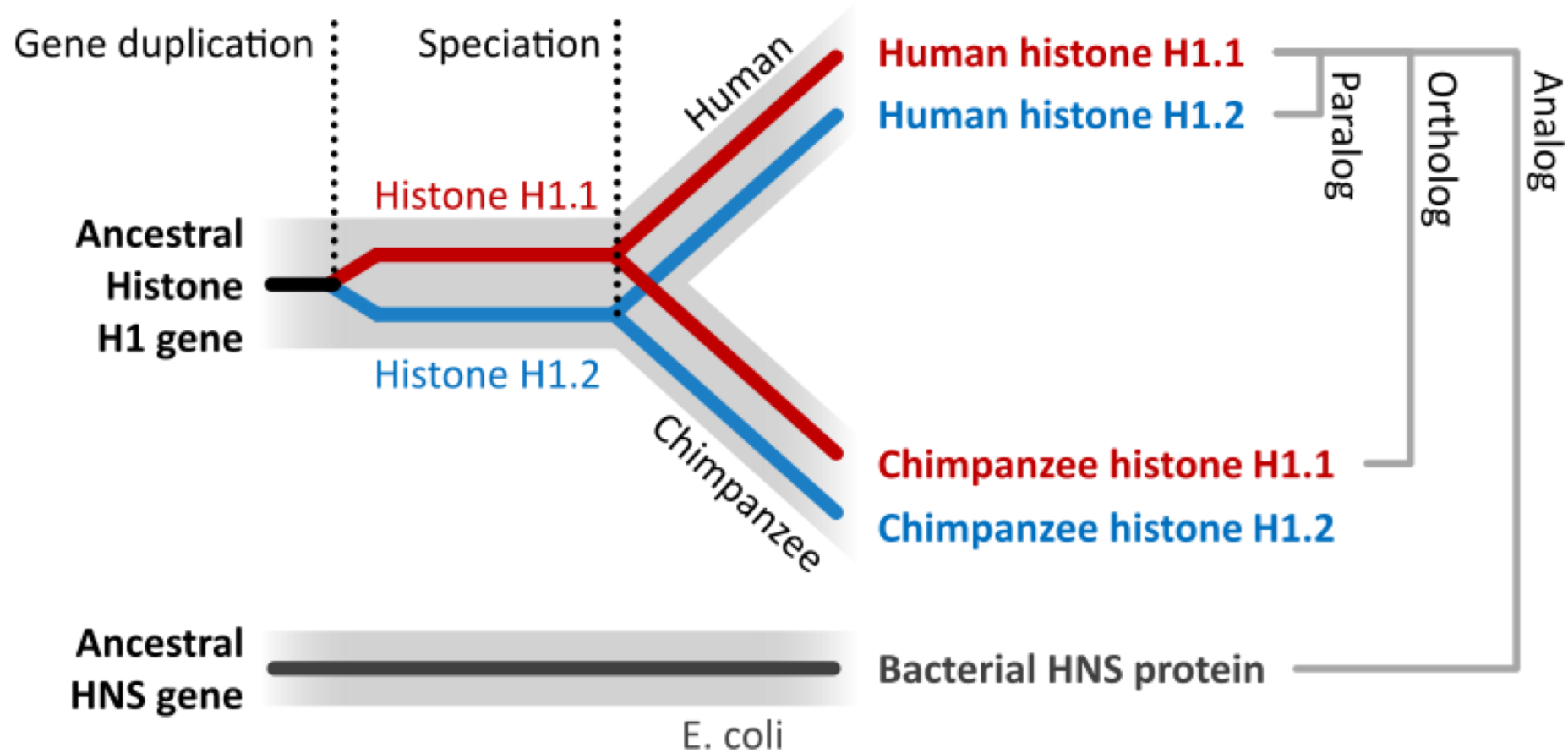


Homology map of human chr7

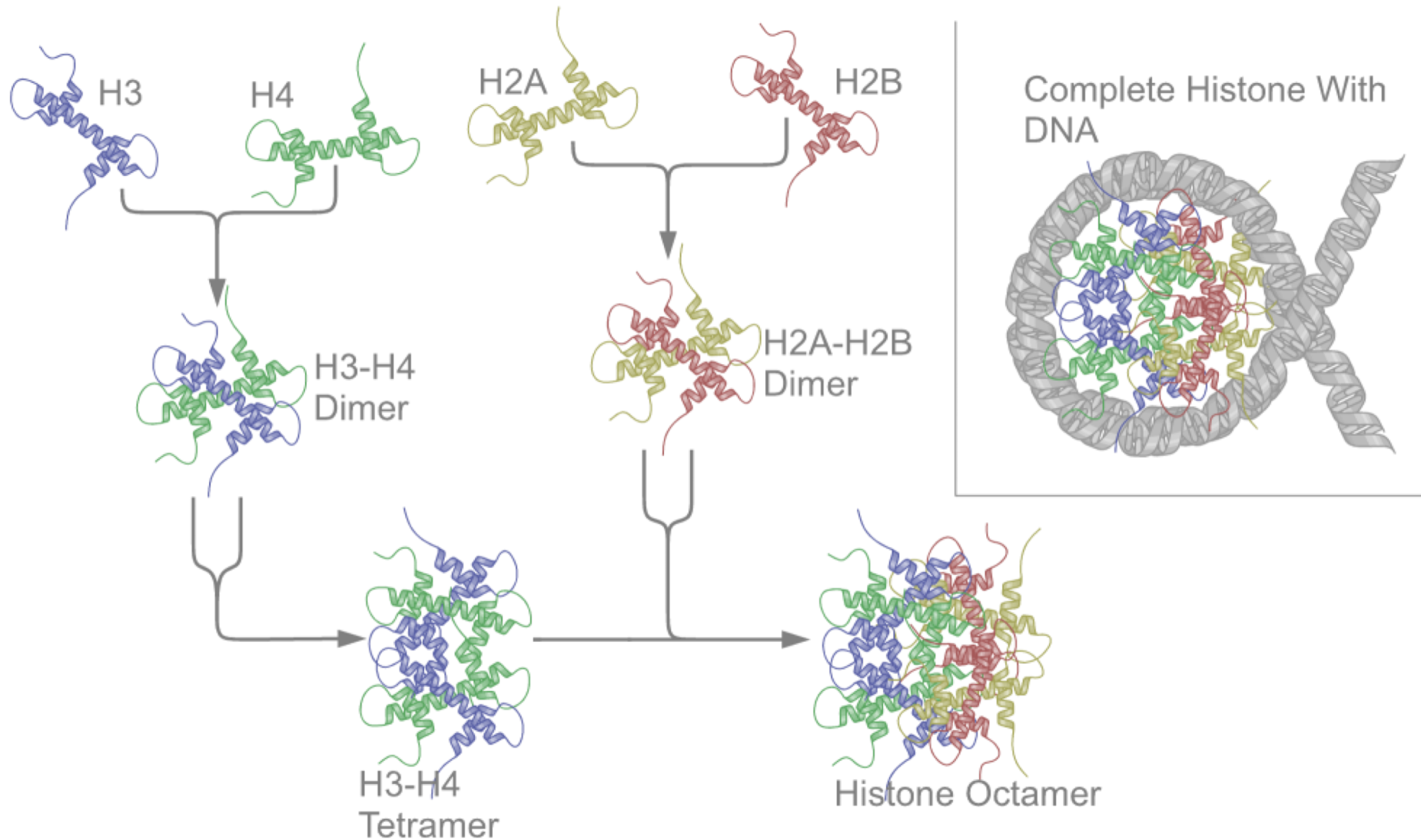


- Homology
 - Shared ancestry
- Showing:
 - >85% identity
 - Blue: 1–5 kbp
 - Red: 5–10 kbp
 - Black: >10 kbp

Homologs: orthologs vs. paralog



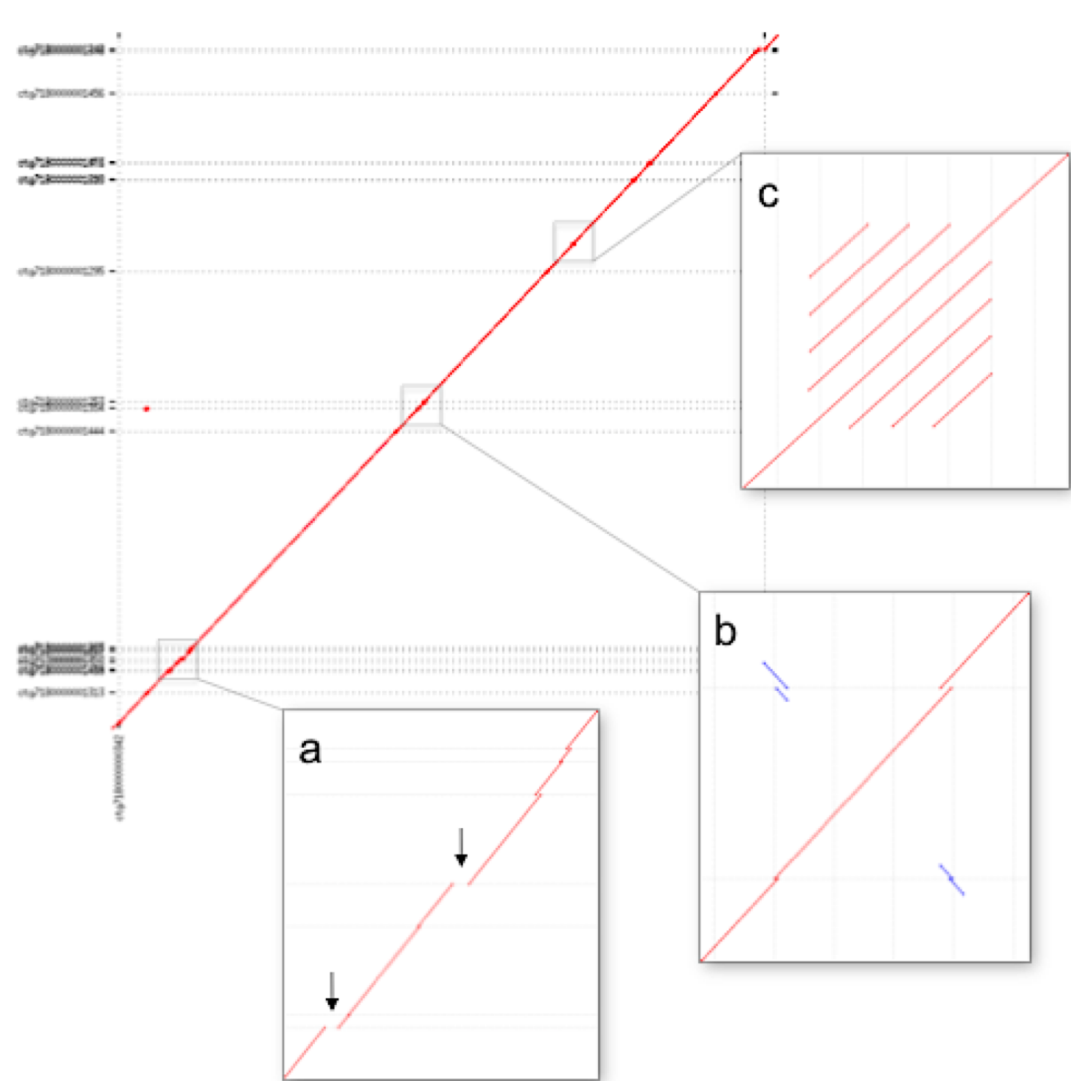
Paralogous genes



Goal of whole genome alignment

- For genomes A and B , find a mapping from each position in A to its orthologous position in B
- Requires an evolutionary model
 - Compared to nucleotide substitution models genome rearrangement models are not well defined
 - Typically based on “synteny” arguments
 - Genomic regions tend to be inherited in blocks

Dotplot quiz



Insertion into Reference R: AIB Q: AB	Insertion into Query R: AB Q: AIB
Collapse Query R: ARRB Q: ARB	Collapse Reference R: ARB Q: ARRB
Collapse Query w/ Insertion R: ARIRB Q: ARB Exact tandem alignment if I=R	Collapse Reference w/ Insertion R: ARB Q: ARIRB Exact tandem alignment if I=R
Collapse Query R: ARRRB Q: ARRB	Collapse Reference R: ARRB Q: ARRRB
Inversion R: ABC Q: AB'C	Rearrangement w/ Disagreement R: ABCDE Q: AFCBE

MUMmer

- Maximal Unique Matcher (MUM)
 - match
 - exact match of a minimum length
 - maximal
 - cannot be extended in either direction without a mismatch
 - unique
 - occurs only once in both sequences (MUM)
 - occurs only once in the reference sequence (MAM)
 - occurs one or more times in either sequence (MEM)

MEMs and MUMs

MUM : maximal unique match

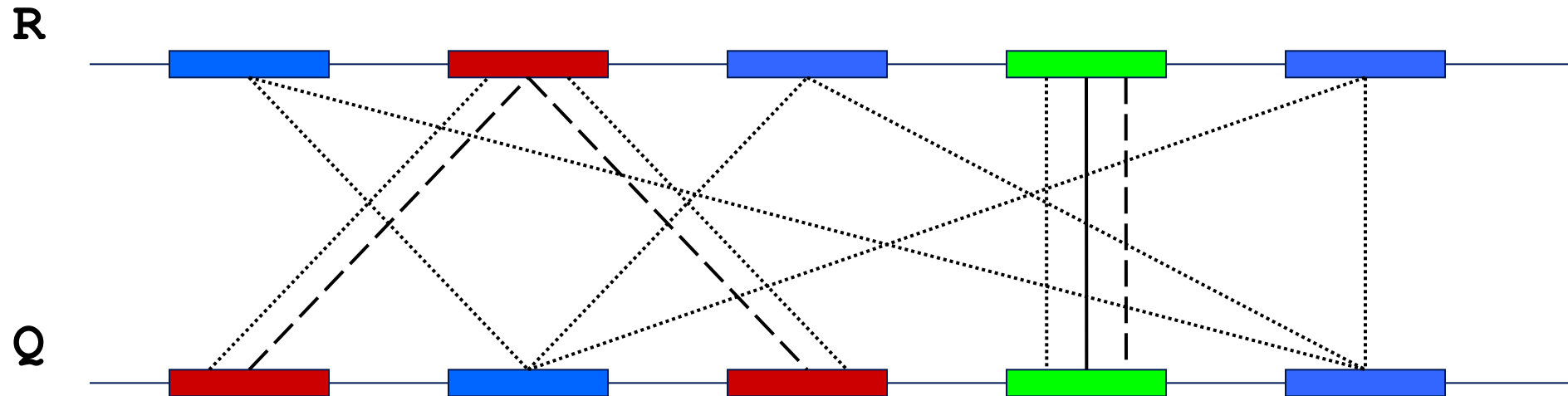
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MAM : maximal almost-unique match

- - - - -

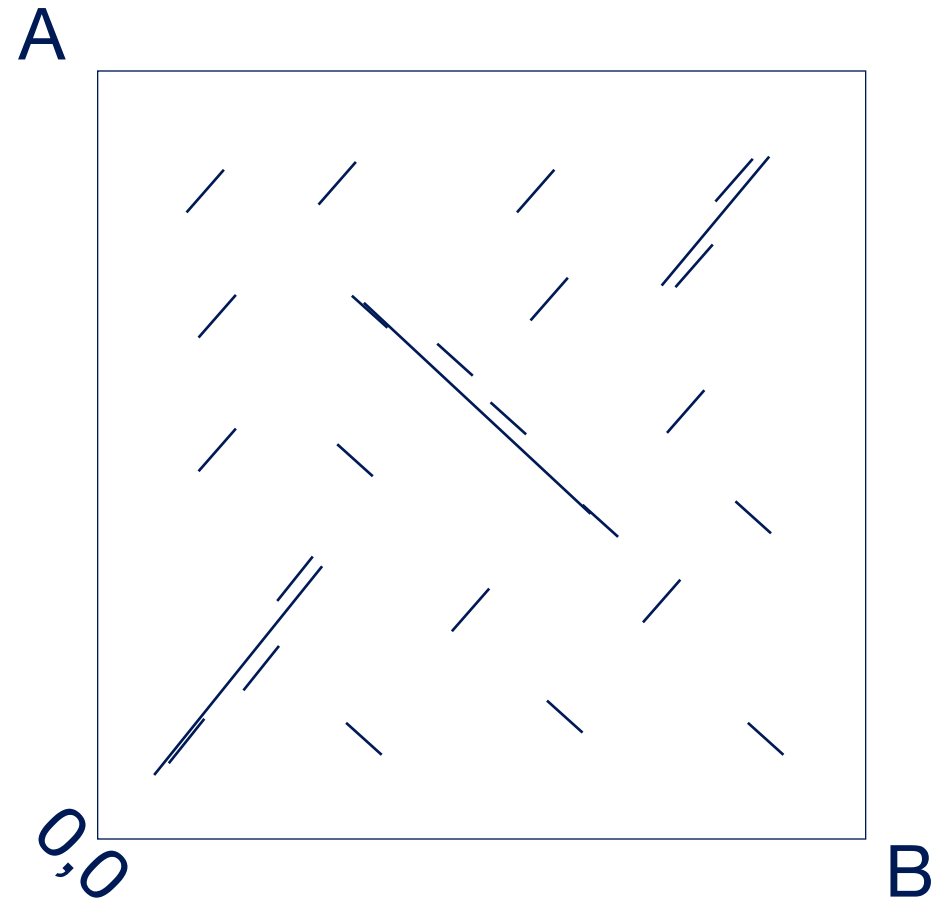
MEM : maximal exact match

.....



How do we turn MUMs into alignments?

- Nucmer approach
 - Find all exact matches
 - Chain exact matches
 - Gapped alignment
- With which algorithms?

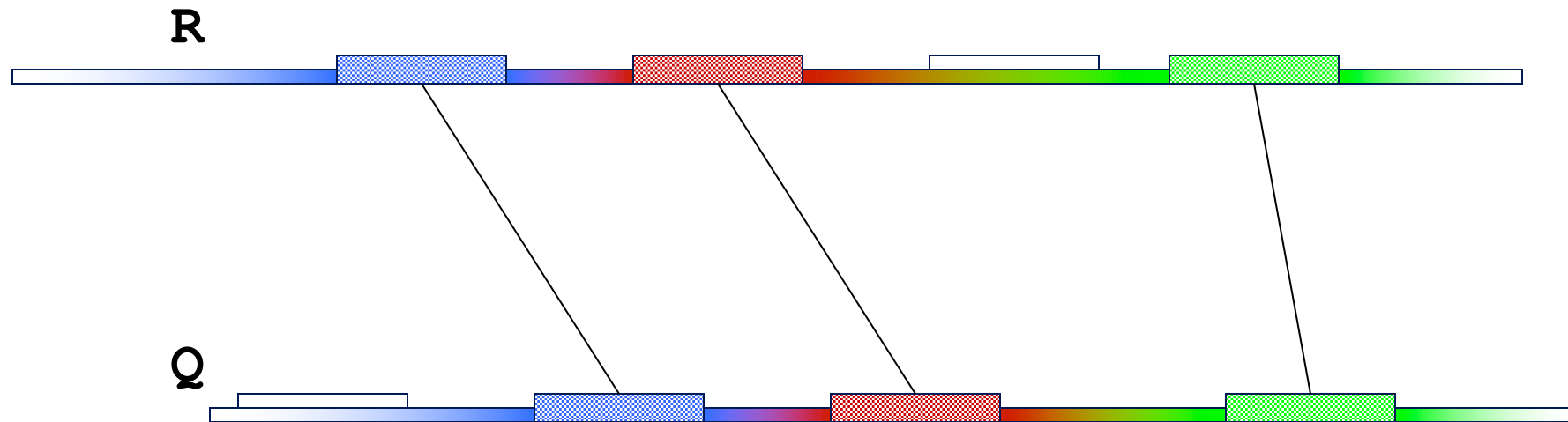


Seed, chain, and extend

FIND all MUMs

CHAIN consistent MUMs

EXTEND alignments



Seed

Suffix indexes for MUM finding

	Suffix tree	Suffix array	FM Index
Time: Does P occur?	$O(n)$	$O(n \log m)$	$O(n)$
Time: Count k occurrences of P	$O(n + k)$	$O(n \log m)$	$O(n)$
Time: Report k locations of P	$O(n + k)$	$O(n \log m + k)$	$O(n + k)$
Space	$O(m)$	$O(m)$	$O(m)$
Needs T?	yes	yes	no
Bytes per input character	>15	~ 4	~ 0.5

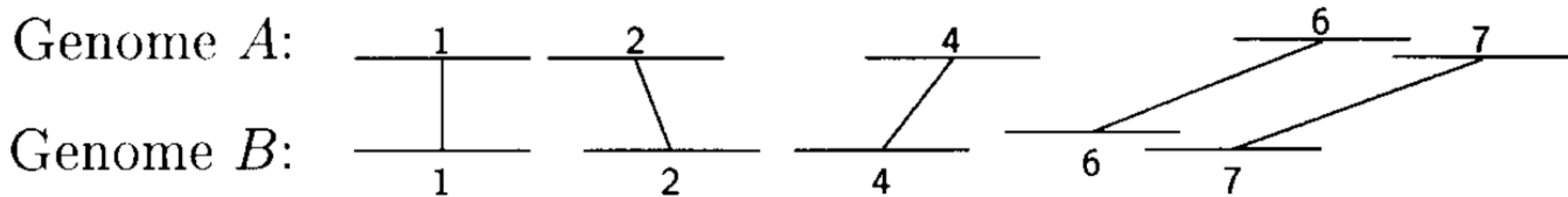
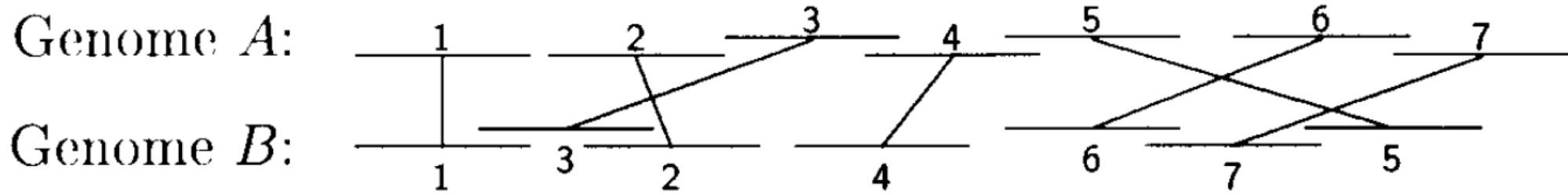
$m = |T|, n = |P|, k = \# \text{ occurrences of } P \text{ in } T$

Chain

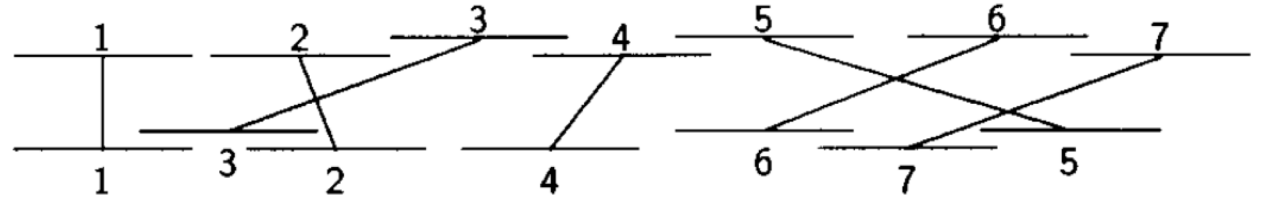
Global: Longest Increasing Subsequence

- A **subsequence** of a **permutation** is a collection of elements of the permutation in the order that they appear. For example, (5, 3, 4) is a subsequence of (5, 1, 3, 4, 2).
- A subsequence is **increasing** if the elements of the subsequence increase. For example, given the permutation (8, 2, 1, 6, 5, 7, 4, 3, 9), an increasing subsequence is (2, 6, 7, 9).
 - **Given:** A permutation π of length n .
 - **Return:** A longest increasing subsequence of π .

Applied LIS example



LIS pseudo code



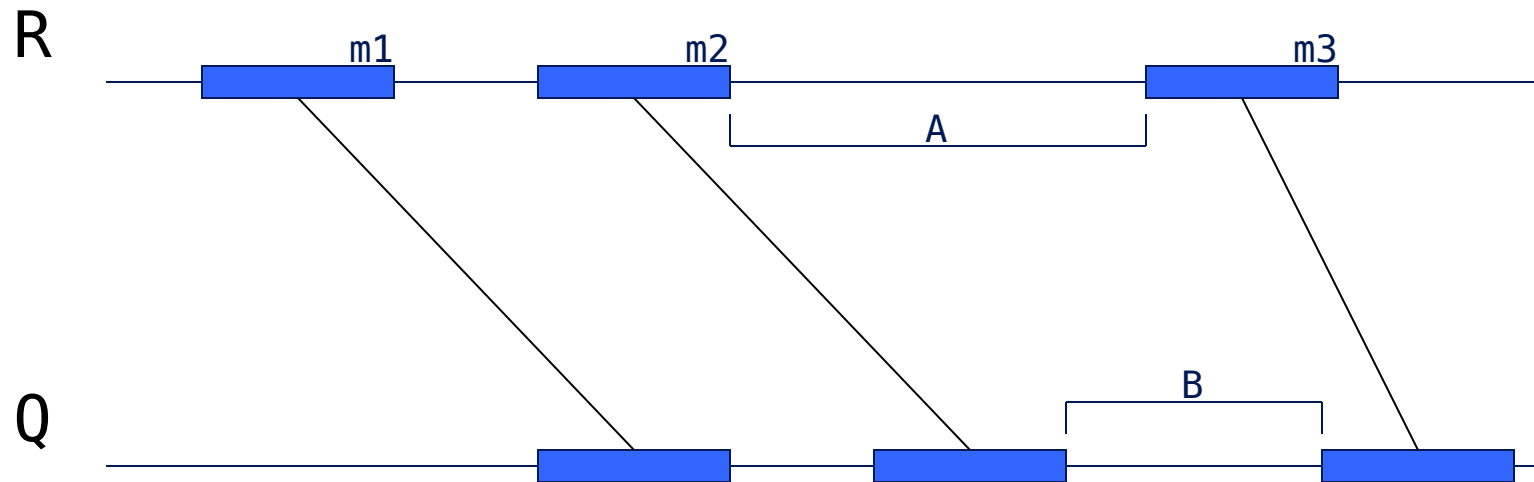
```
for i in [1..n]
  for j in [1..i)
    score = max[j] + len(i) - overlap(j,i)
    if score > max[i]
      max[i] = score
      pre[i] = j
```

Local: Agglomerative Clustering

cluster length = $\sum m_i$

gap distance = A

diagonal difference = $|B - A|$ or $|B - A| / B$



Agglomerative pseudo code

```
for i in [1..n)
  for j in [i+1..n]
    if gap(i,j) < g and diagdiff(i,j) < d
      union(find(i), find(j))
```



Extend



T
A

T
A
G
G

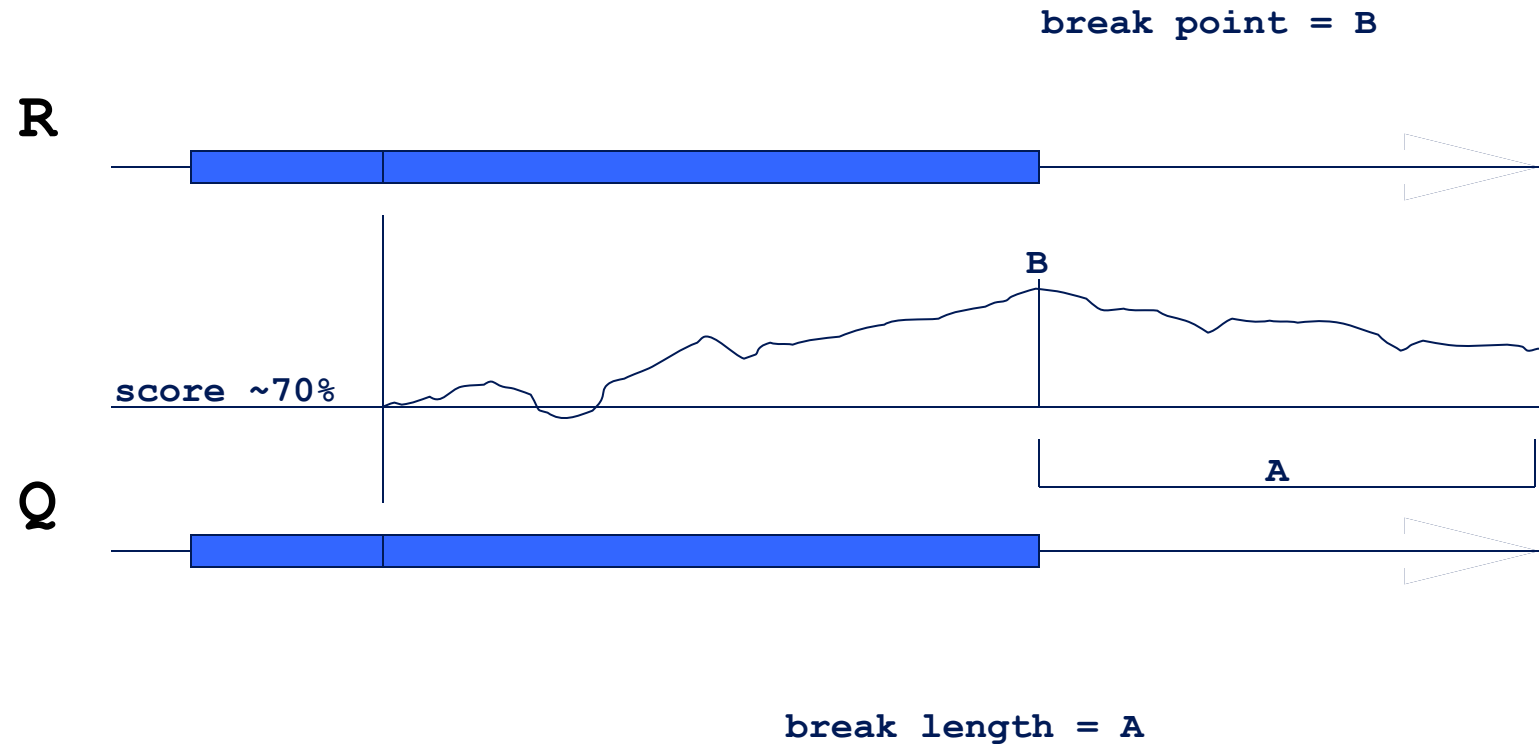
Banded alignment extension

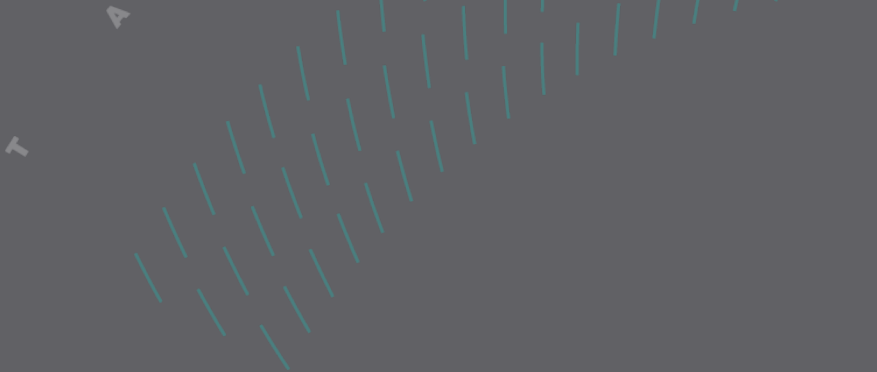
B

A

	^	T	T	G	C	A	G
^	0 ↓	1 ↘	2 →	3* →	4 →	5 →	6 →
T	1 ↓	0 ↓	1 ↘	2 →	3 →	4 →	5 ↘
G	2 ↓	1 ↓	1 ↘	1 ↓	2 →	3 →	4 ↘
C	3* ↓	2 ↓	2 ↓	2 ↓	1 ↘	2 →	3 →
T	4 ↓	3 ↓	2 ↘	3* ↓	2 ↓	2 ↘	3* ↘
G	5 ↓	4 ↓	3 ↓	2 ↘	3 ↓	3* ↓	2 ↘

Alignment extension

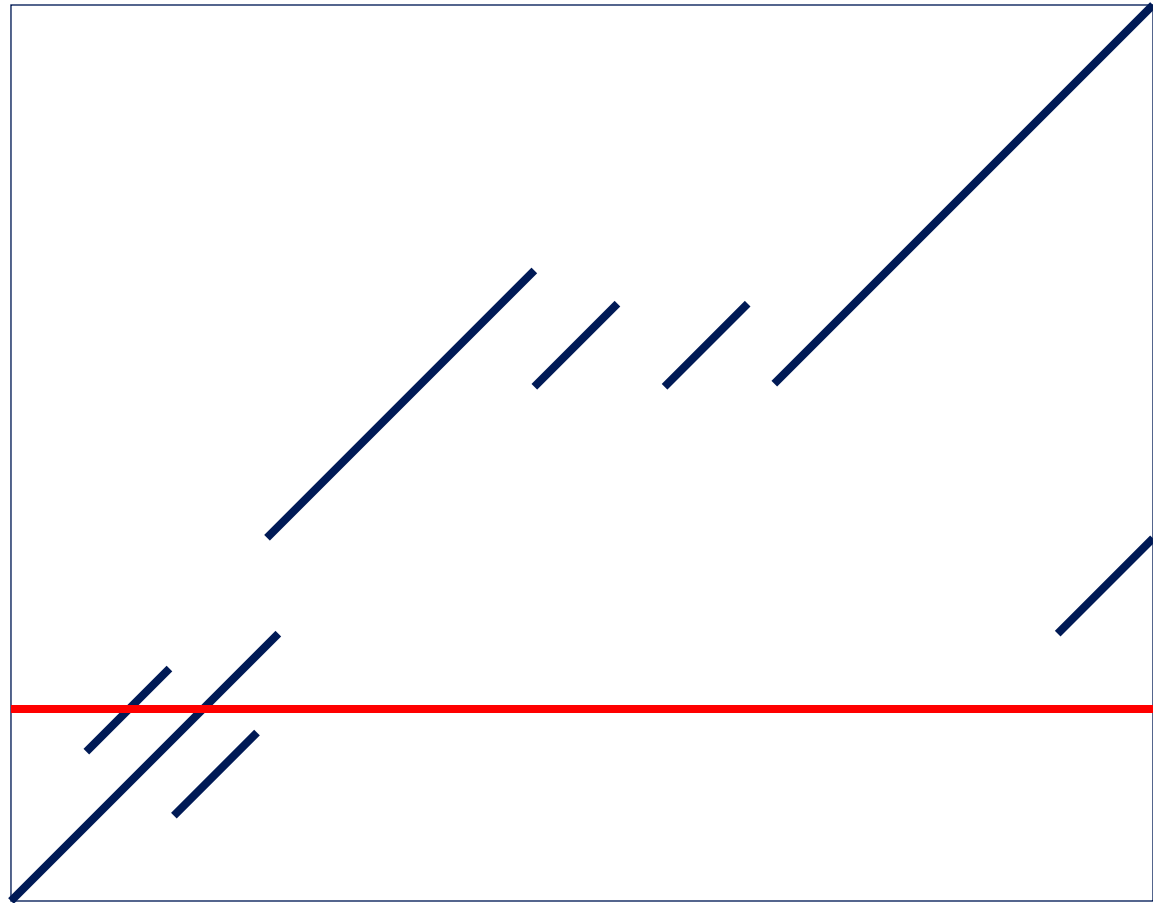




Filter

Which are orthologs?

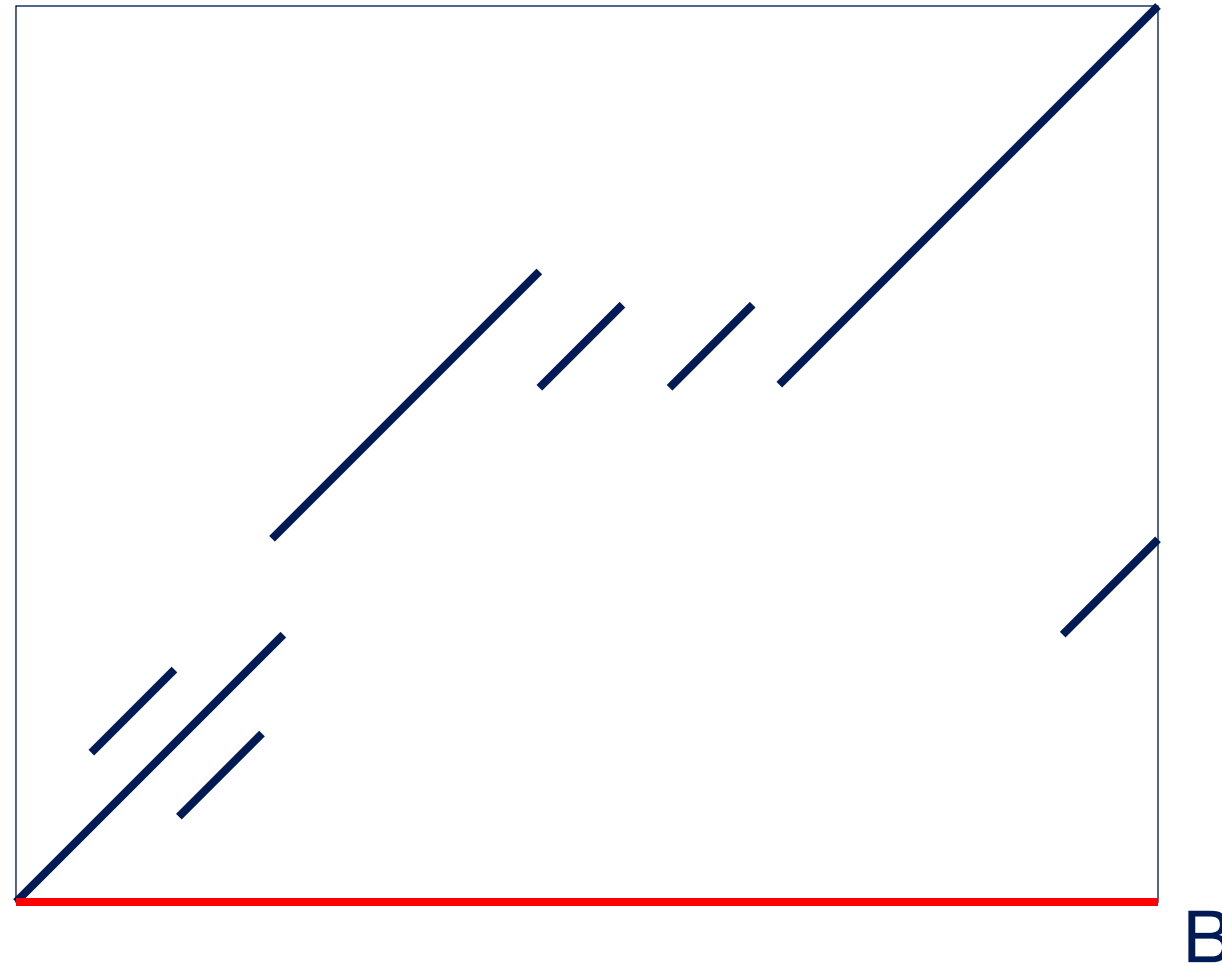
A



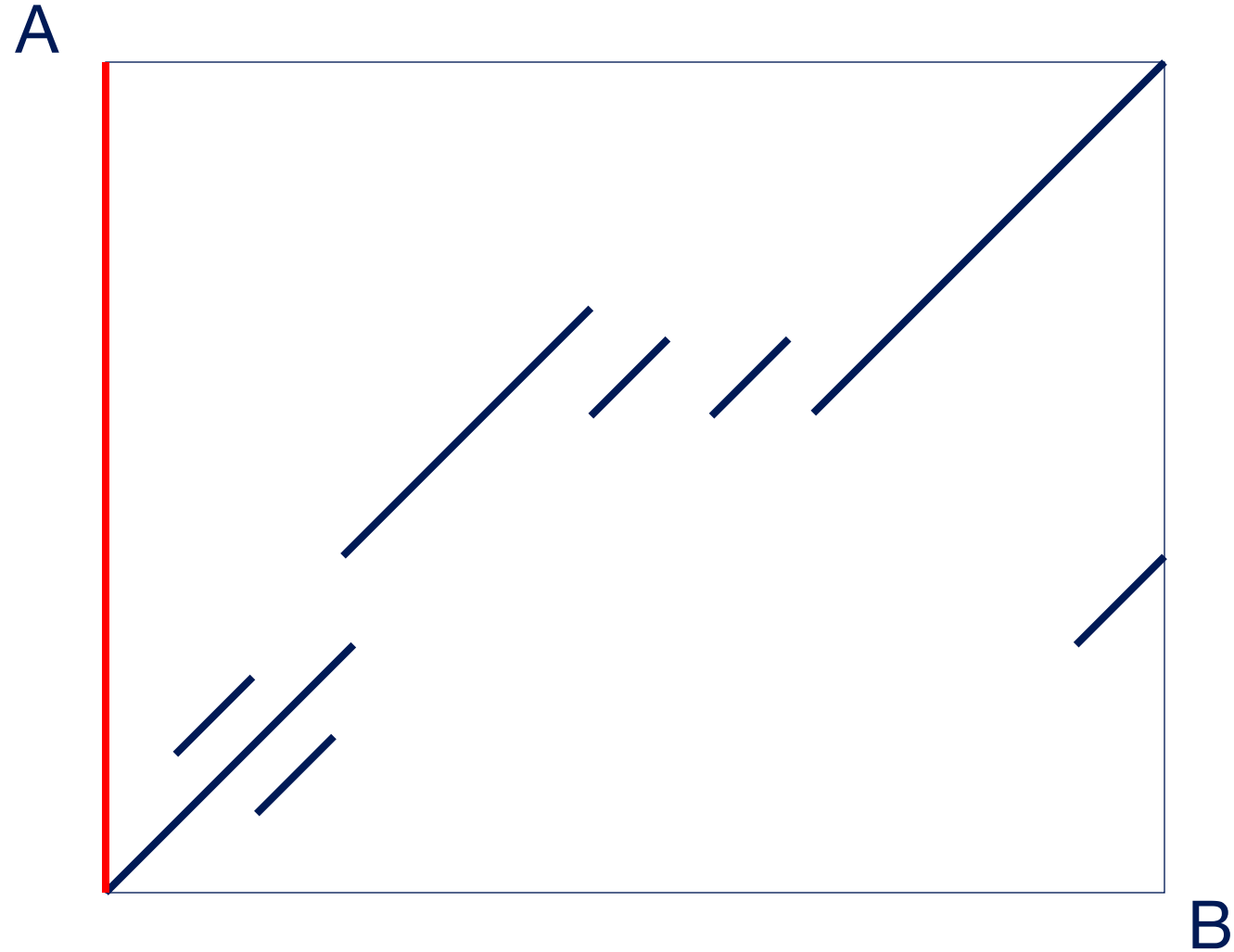
B

Best for A

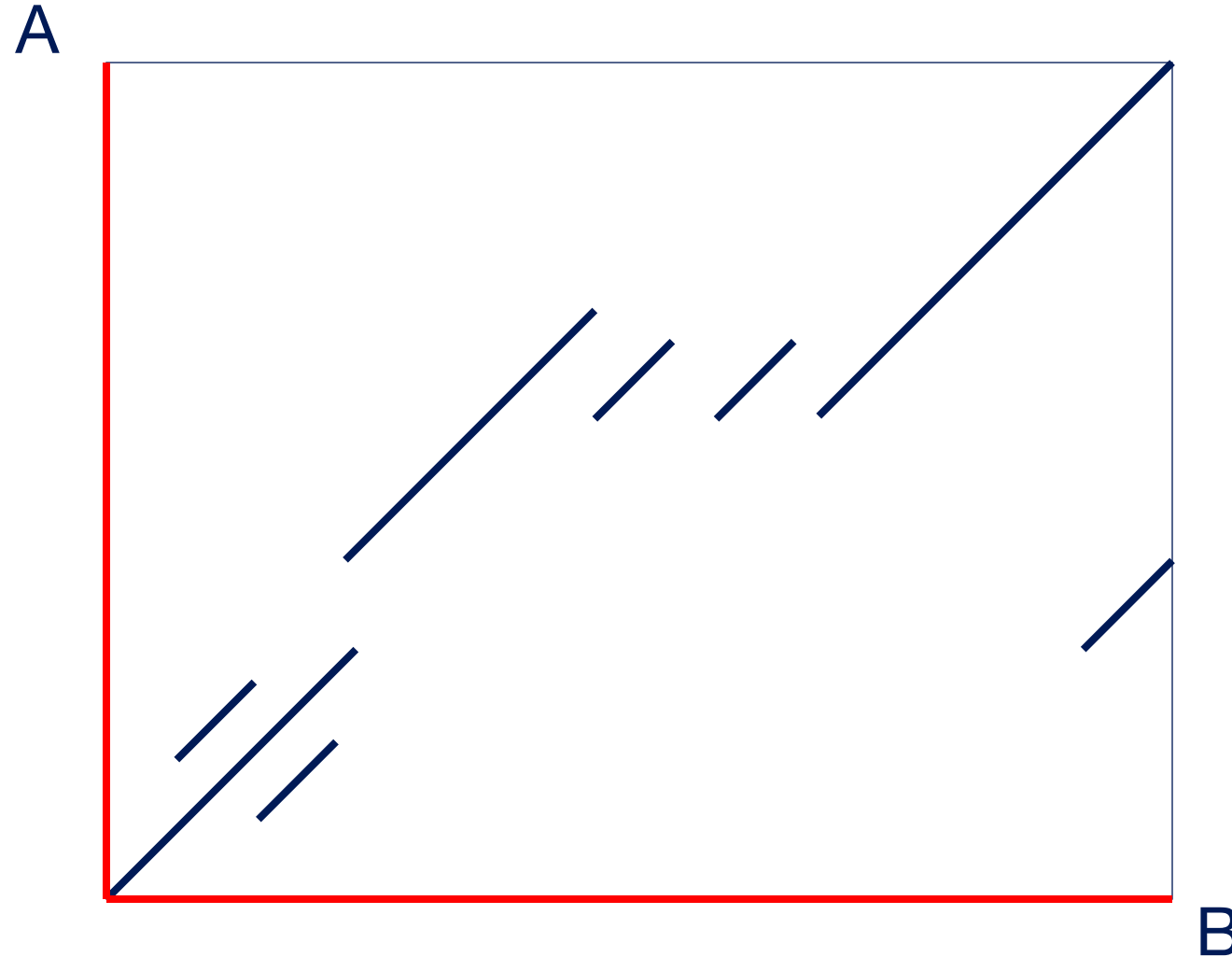
A



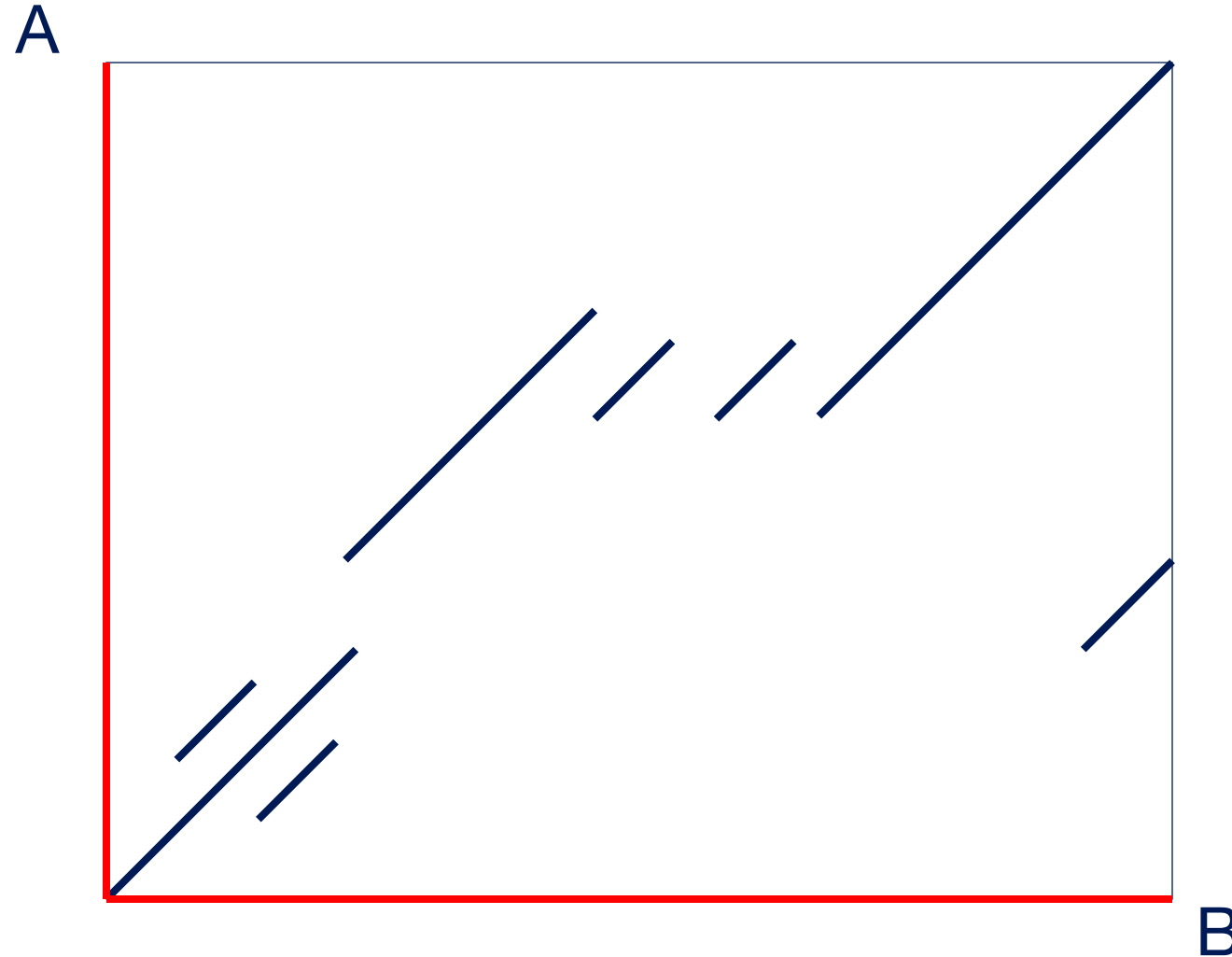
Best for B



Best for *A and B* (orthologs only)



Best for *A* or *B* (paralogs too)



SuperMap

