Reasoning about Knowledge and Beliefs Lecture 9

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October 9, 2013

Conceptions of Belief

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Eric Schwitzgebel. Belief. In The Stanford Encyclopedia of Philosophy.

Franz Huber. Formal Theories of Belief. In The Stanford Encyclopedia of Philosophy.

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- D. Christensen. Putting Logic in its Place. Oxford University Press.
- H. Leitgeb. The Lockean Thesis Revisited. Working Paper, 2010.

Epistemic-Probability Models

Adding Probabilities

Epistemic-Probability Model: $\mathcal{M} = \langle W, \{\sim_i\}_{i \in \mathcal{A}}, \{P_i\}_{i \in \mathcal{A}}, V \rangle$ where each \sim_i is an equivalence relation on W is an epistemic model and $P_i : W \to \Delta(W)$ assigns to each state a probability measure over W, and V is a valuation function.

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Write p_i^w for the *i*'s probability measure at state *w*. We make two natural assumptions:

1. For all
$$v \in W$$
, if $p_i^w(v) > 0$ then $p_i^w = p_i^v$; and
2. For all $v \notin [w]_i$, $p_i^w(v) = 0$.

Common Prior

Epistemic-Probabilistic Models: $\mathcal{M} = \langle W, \{\sim_i\}_{i \in \mathcal{A}}, \mathbf{p}, V \rangle$

Common Prior: $p: W \rightarrow [0, 1]$ is a probability measure (assume *W* finite)

Truth: $\mathcal{M}, w \models \varphi$ is defined as follows:

•
$$\mathcal{M}, w \models p \text{ iff } w \in V(p) \text{ (with } p \in At)$$

•
$$\mathcal{M}, w \models \neg \varphi$$
 if $\mathcal{M}, w \not\models \varphi$

$$\blacktriangleright \ \mathcal{M}, w \models \varphi \land \psi \text{ if } \mathcal{M}, w \models \varphi \text{ and } \mathcal{M}, w \models \psi$$

- $\mathcal{M}, w \models K_i \varphi$ if for each $v \in W$, if $w \sim_i v$, then $\mathcal{M}, v \models \varphi$
- $\blacktriangleright \mathcal{M}, w \models B^{r} \varphi \text{ iff } p(\llbracket \varphi \rrbracket \mid \llbracket w \rrbracket_{i}) = \frac{p(\llbracket \varphi \rrbracket \cap [w]_{i})}{p(\llbracket w \rrbracket_{i})} \ge r$

$$B_i^0 \varphi \\ B_i^1 \top$$

$$\blacktriangleright B_i^q(\varphi \land \psi) \land B_i^p(\varphi \land \neg \psi) \to B_i^{q+p}\varphi, \quad q+p \leq 1$$

$$\neg B_i^q(\varphi \land \psi) \land \neg B_i^p(\varphi \land \neg \psi) \to \neg B_i^{q+p}\varphi, \quad q+p \le 1$$

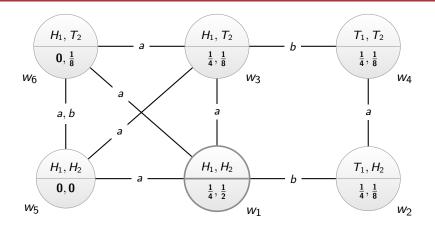
$$\blacktriangleright B_i^q \varphi \to \neg B_i^p \neg \varphi, \quad q+p>1$$

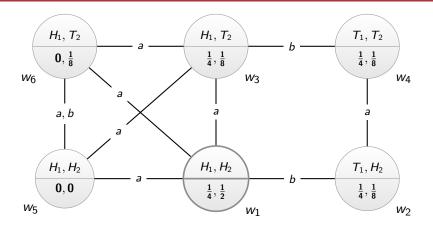
Archimedian Rule: If $\psi \to B_i^p \varphi$ is valid for each p < q, then $\psi \to B_i^q \varphi$ is valid.

- $K_i \varphi \to B_i^q \varphi$
- $\blacktriangleright B_i^q \varphi \to K_i B_i^q \varphi$
- $\blacktriangleright \neg B_i^q \varphi \to K_i \neg B_i^q \varphi$
- if $\varphi \to \psi$ is valid then so is $B_i^q \varphi \to B_i^q \psi$

- $K_i \varphi \to B_i^q \varphi$ $B_i^q \varphi \to K_i B_i^q \varphi$ $\neg B_i^q \varphi \to K_i \neg B_i^q \varphi$
- if $\varphi \to \psi$ is valid then so is $B_i^q \varphi \to B_i^q \psi$

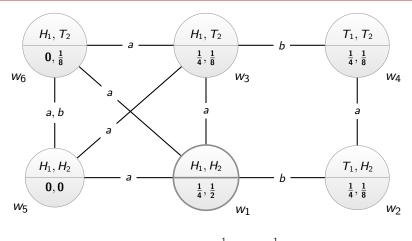
$\pi_i(\llbracket \varphi \rrbracket_{\mathcal{M}} \mid \llbracket B_i^q \varphi \rrbracket_{\mathcal{M}}) \geq q$



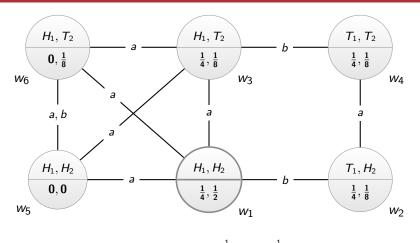


 $\blacktriangleright \mathcal{M}, w_1 \models \neg K_a H_2 \land \neg K_a T_2 \land B_a^{\frac{1}{2}} H_2 \land B_a^{\frac{1}{2}} T_2$

Reasoning about Knowledge and Beliefs



$$\mathcal{M}, w_1 \models \neg K_a H_2 \land \neg K_a T_2 \land B_a^{\frac{1}{2}} H_2 \land B_a^{\frac{1}{2}} T_2$$
$$\mathcal{M}, w_1 \models \neg K_b H_1 \land \neg K_b T_1 \land B_b^{\frac{4}{5}} H_1 \land B_b^{\frac{1}{5}} T_1$$



$$\mathcal{M}, w_1 \models \neg K_a H_2 \land \neg K_a T_2 \land B_a^{\frac{1}{2}} H_2 \land B_a^{\frac{1}{2}} T_2 \mathcal{M}, w_1 \models \neg K_b H_1 \land \neg K_b T_1 \land B_b^{\frac{4}{5}} H_1 \land B_b^{\frac{1}{5}} T_1 \mathcal{M}, w_1 \models \neg K_a (K_b H_2 \lor K_b T_2) \land B_a^{1} (K_b H_2 \lor K_b T_2)$$