

# CMSC 132:

# Object-Oriented Programming II

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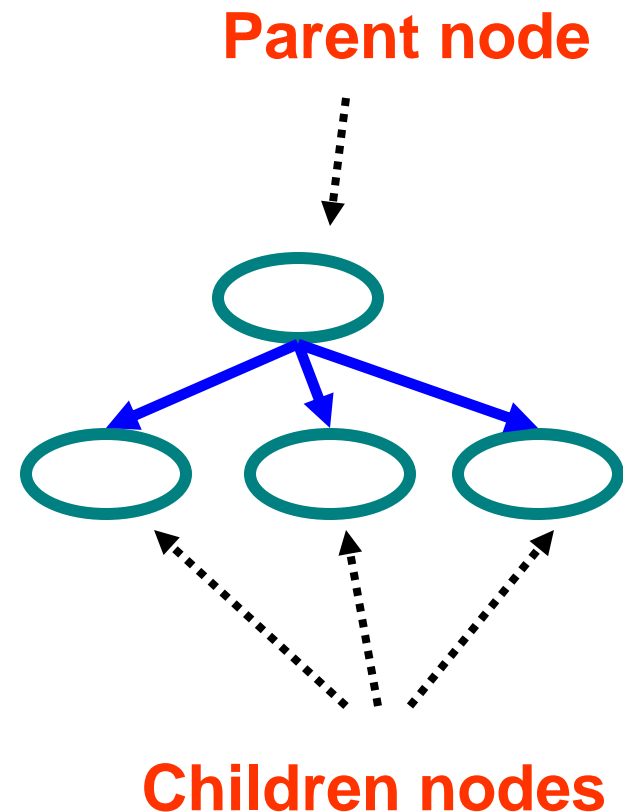


## Trees & Binary Search Trees

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University of Maryland, College Park

# Trees

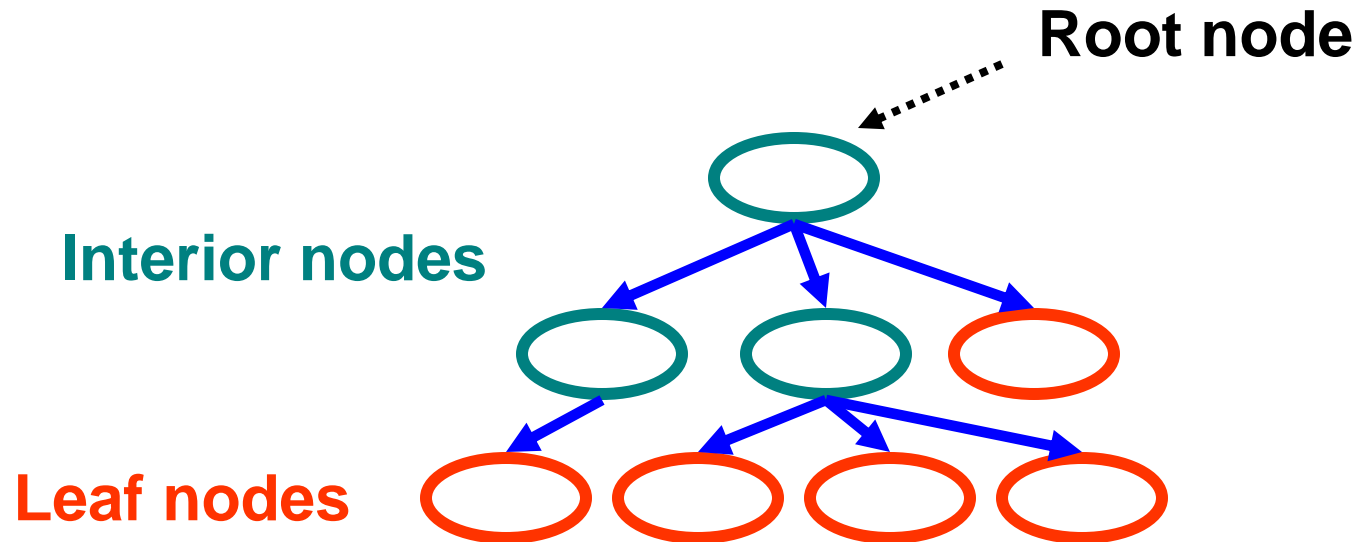
- Trees are hierarchical data structures
  - One-to-many relationship between elements
- Tree node / element
  - Contains data
  - Referred to by only 1 (**parent**) node
  - Contains links to any number of (**children**) nodes



# Trees

## ■ Terminology

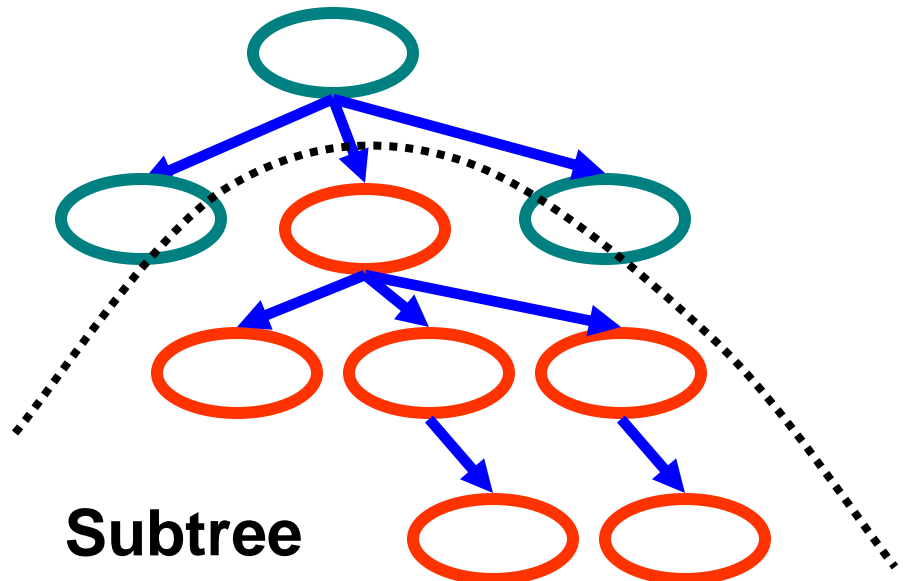
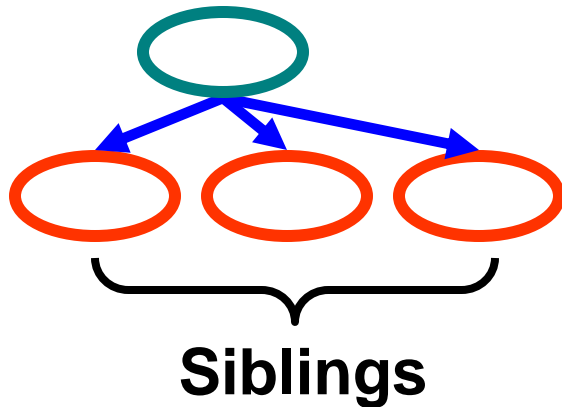
- Root  $\Rightarrow$  node with no parent
- Leaf  $\Rightarrow$  all nodes with no children
- Interior  $\Rightarrow$  all nodes with children



# Trees

## ■ Terminology

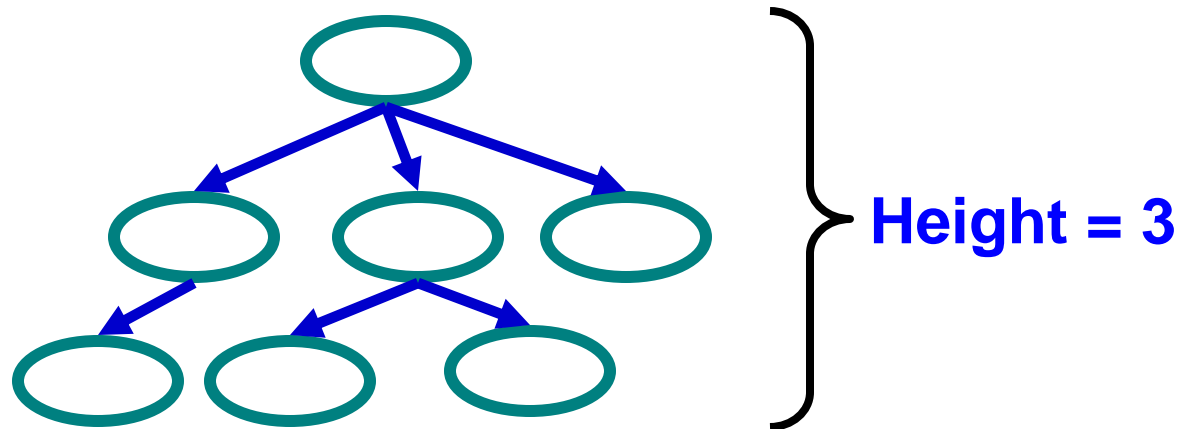
- Sibling  $\Rightarrow$  node with same parent
- Descendent  $\Rightarrow$  children nodes & their descendants
- Subtree  $\Rightarrow$  portion of tree that is a tree by itself  
 $\Rightarrow$  a node and its descendants



# Trees

## ■ Terminology

- **Level**  $\Rightarrow$  is a measure of a node's distance from root
- **Definition of level**
  - If node is the root of the tree, its level is 1
  - Else, the node's level is  $1 +$  its parent's level
- **Height (depth)**  $\Rightarrow$  max level of any node in tree

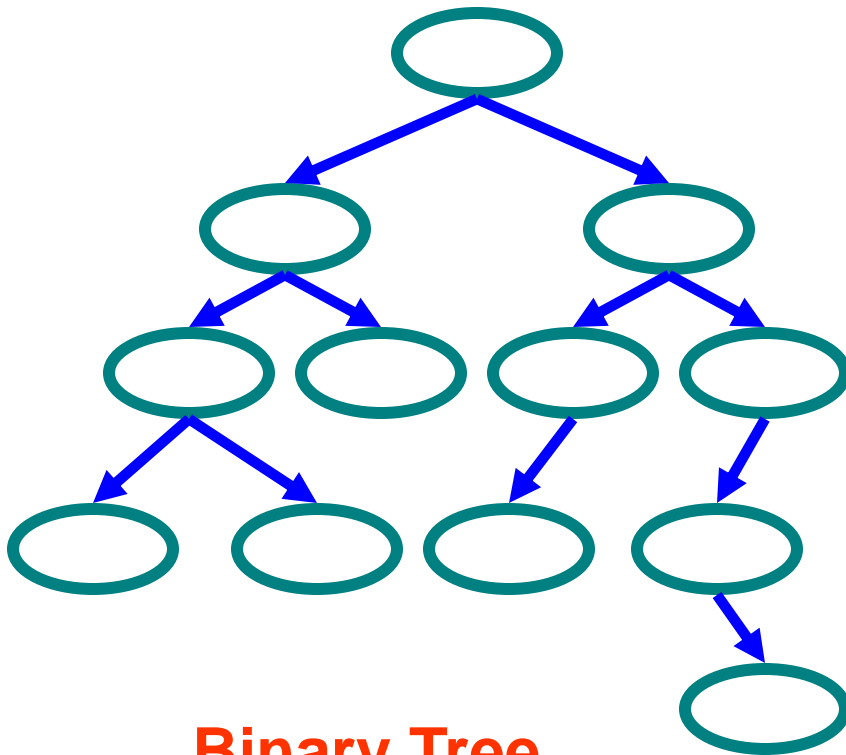


# Binary Trees

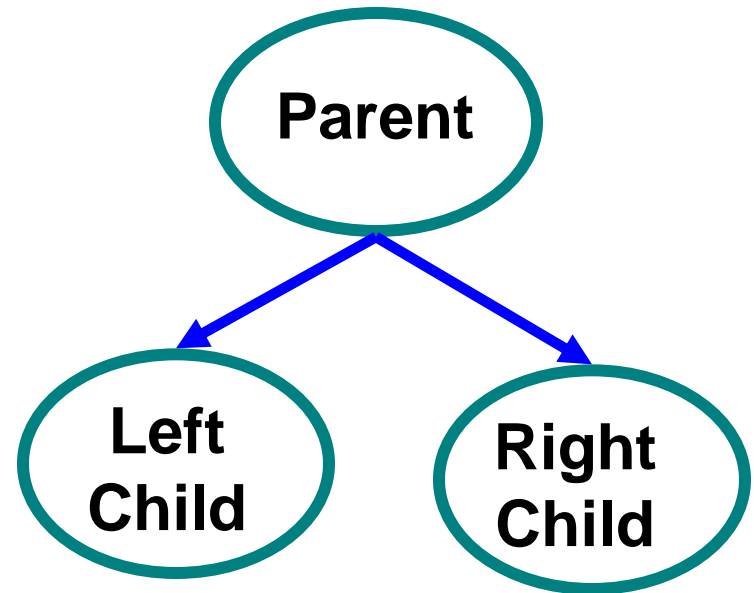
## ■ Binary tree

■ Tree with 0–2 children per node

■ Left & right child / subtree



**Binary Tree**



# Tree Traversal

- Often we want to

1. Find all nodes in tree
2. Determine their relationship

- Can do this by

1. Walking through the tree in a prescribed order
2. Visiting the nodes as they are encountered

- Process is called **tree traversal**

# Tree Traversal

## ■ Goal

- Visit every node in binary tree

## ■ Approaches

### ■ Depth first

- Preorder  $\Rightarrow$  parent before children
- Inorder  $\Rightarrow$  left child, parent, right child
- Postorder  $\Rightarrow$  children before parent

- Breadth first  $\Rightarrow$  closer nodes first



# Tree Traversal Methods

## ■ Pre-order

1. Visit **node** // first
2. Recursively visit left subtree
3. Recursively visit right subtree

## ■ In-order

1. Recursively visit left subtree
2. Visit **node** // second
3. Recursively right subtree

## ■ Post-order

1. Recursively visit left subtree
2. Recursively visit right subtree
3. Visit **node** // last

# Tree Traversal Methods

## ■ Breadth-first

```
BFS(Node n) {  
    Queue Q = new Queue();  
    Q.enqueue(n);           // insert node into Q  
    while ( !Q.empty()) {  
        n = Q.dequeue();    // remove next node  
        if ( !n.isEmpty()) {  
            visit(n);        // visit node  
            Q.enqueue(n.Left()); // insert left subtree in Q  
            Q.enqueue(n.Right()); // insert right subtree in Q  
        }  
    }  
}
```

# Tree Traversal Examples

## ■ Pre-order (prefix)

■  $+ \times 2 3 / 8 4$

## ■ In-order (infix)

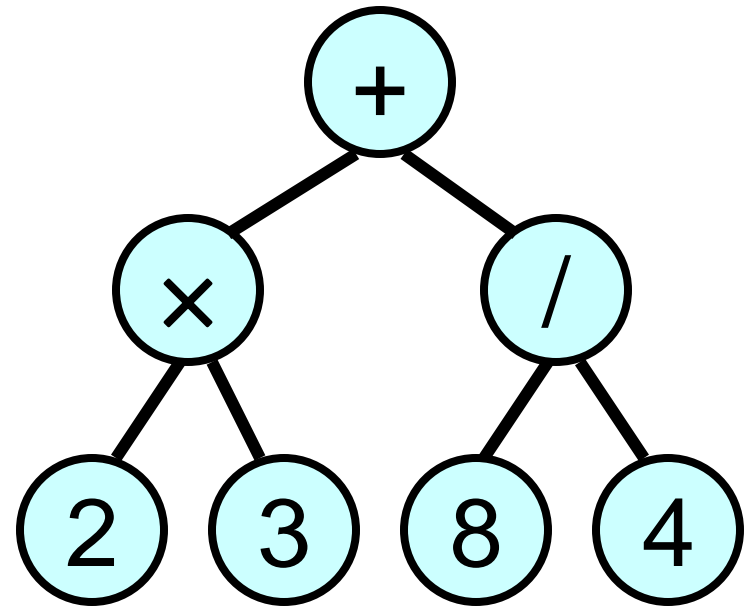
■  $2 \times 3 + 8 / 4$

## ■ Post-order (postfix)

■  $2 3 \times 8 4 / +$

## ■ Breadth-first

■  $+ \times / 2 3 8 4$



Expression tree

# Binary Tree Implementation

- Using a class to represent a Node

```
Class Node {  
    KeyType key;  
    Node left, right;    // null if empty  
}
```

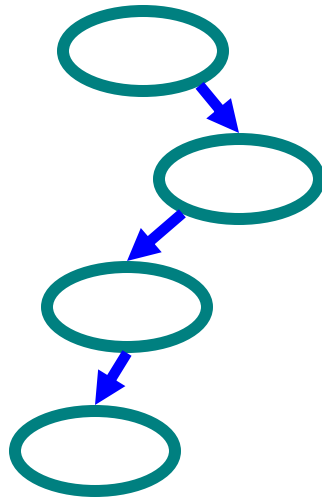
Node root = null; // Empty Tree

- Using a Polymorphic Binary Tree
  - We will talk about this implementation later on

# Types of Binary Trees

## ■ Degenerate

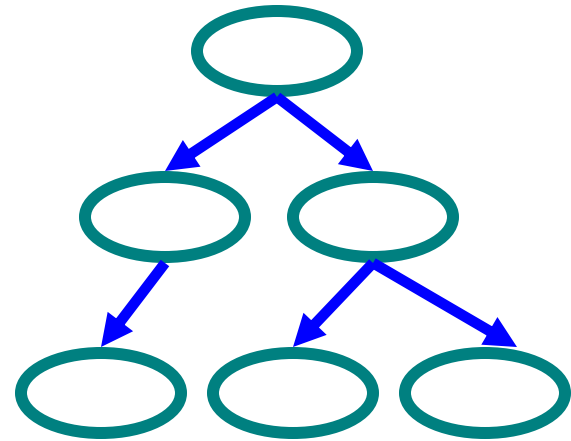
- Mostly 1 child / node
- Height =  $O(n)$
- Similar to linear list



**Degenerate  
binary tree**

## ■ Balanced

- Mostly 2 child / node
- Height =  $O(\log(n))$
- $2^{\text{Height}} - 1 = n$  (# of nodes)
- Useful for searches



**Balanced  
binary tree**

# Binary Search Trees

- **Key property**

- **Value at node**

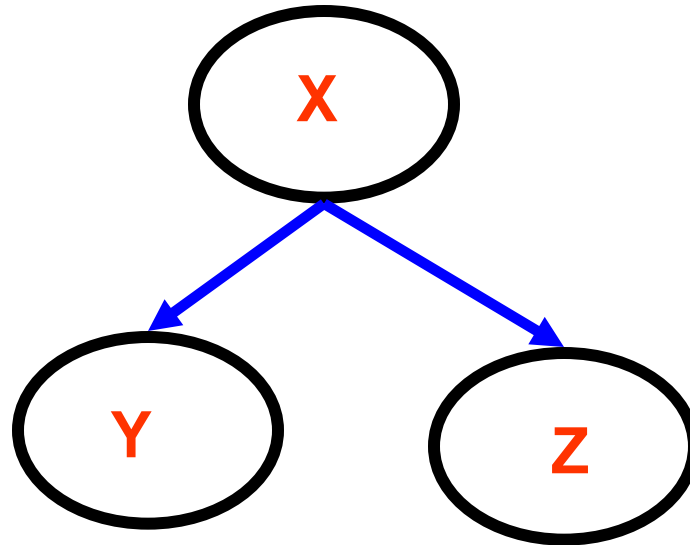
- **Smaller values in left subtree**

- **Larger values in right subtree**

- **Example**

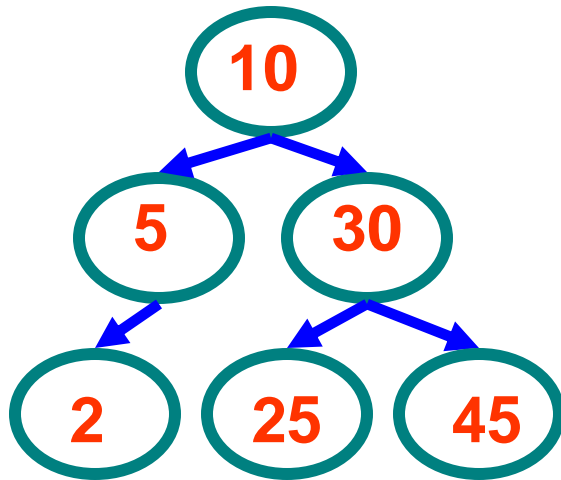
- $X > Y$

- $X < Z$

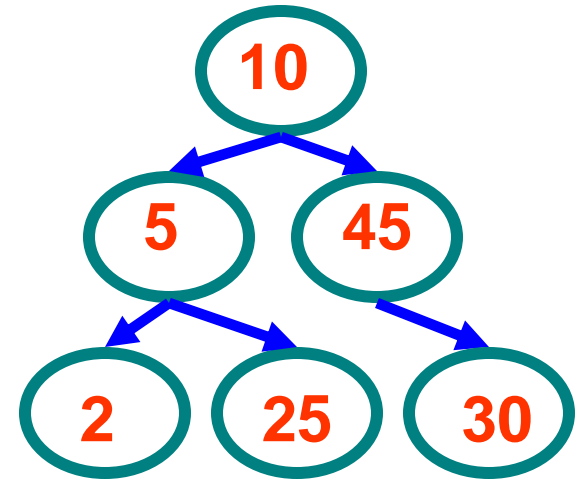
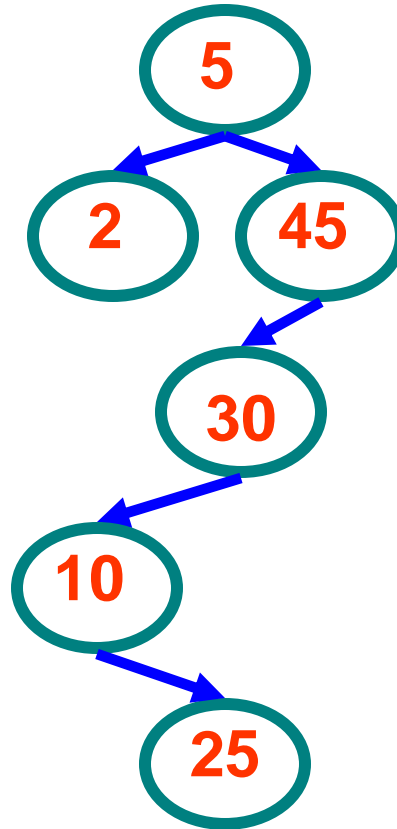


# Binary Search Trees

## ■ Examples



Binary  
search trees



Non-binary  
search tree

# Tree Traversal Examples

## ■ Pre-order

■ 44, 17, 32, 78,  
50, 48, 62, 88

## ■ In-order

■ 17, 32, 44, 48,  
50, 62, 78, 88

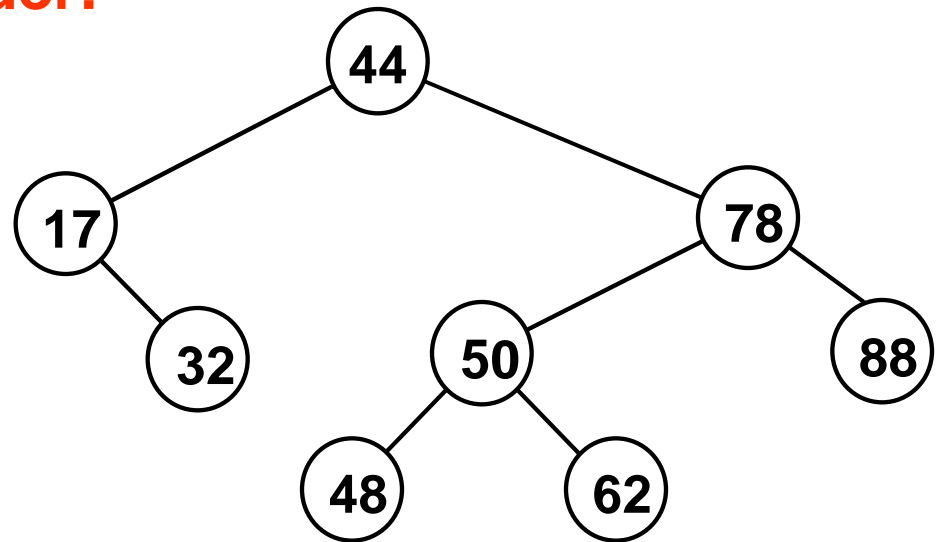
## ■ Post-order

■ 32, 17, 48, 62,  
50, 88, 78, 44

## ■ Breadth-first

■ 44, 17, 78, 32,  
50, 88, 48, 62

Sorted  
order!

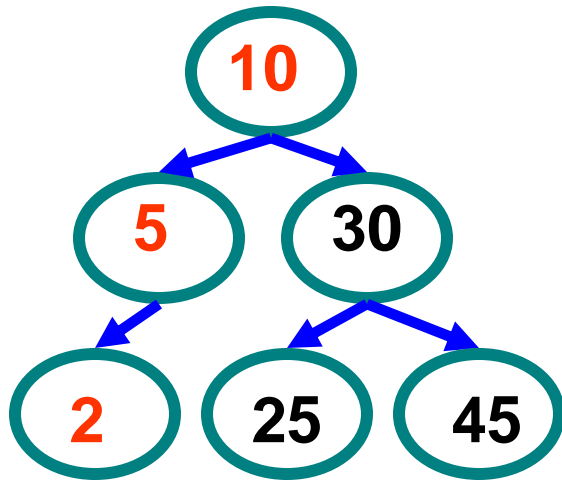


Binary search tree



# Example Binary Searches

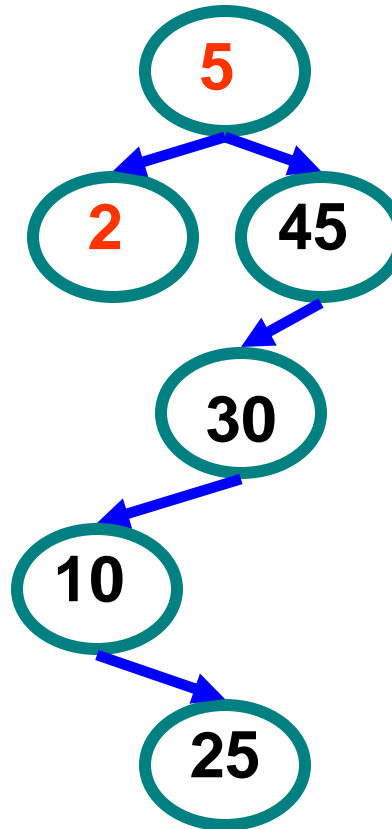
## ■ Find ( 2 )



$2 < 10$ , left

$2 < 5$ , left

$2 = 2$ , found

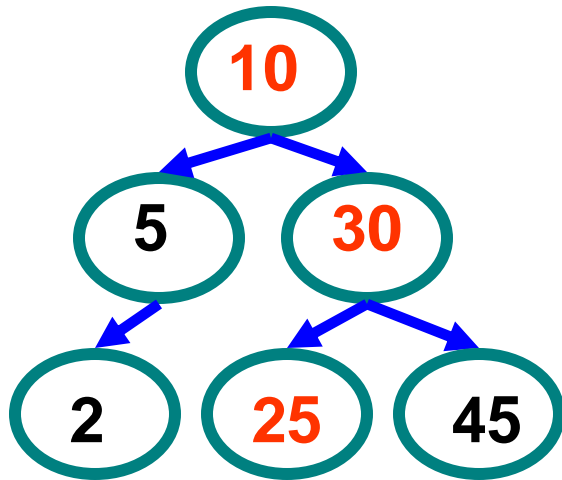


$2 < 5$ , left

$2 = 2$ , found

# Example Binary Searches

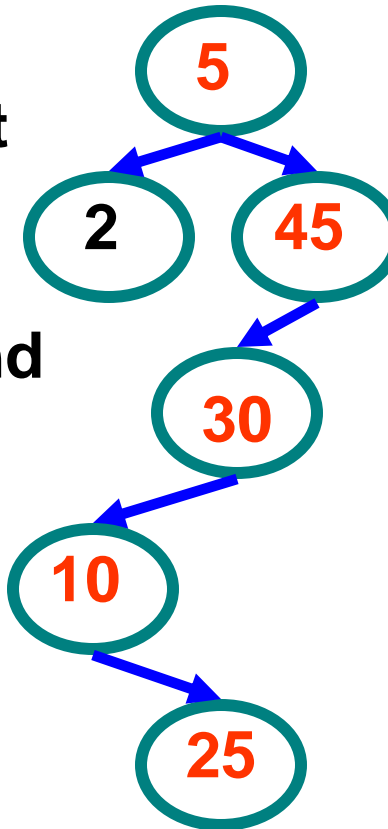
## ■ Find ( 25 )



$25 > 10$ , right

$25 < 30$ , left

$25 = 25$ , found



$25 > 5$ , right

$25 < 45$ , left

$25 < 30$ , left

$25 > 10$ , right

$25 = 25$ , found

# Binary Search Properties

## ■ Time of search

- Proportional to height of tree

- Balanced binary tree

  - $O(\log(n))$  time

- Degenerate tree

  - $O(n)$  time

  - Like searching linked list / unsorted array

## ■ Requires

- Ability to compare key values

# Binary Search Tree Construction

- **How to build & maintain binary trees?**
  - **Insertion**
  - **Deletion**
- **Maintain key property (invariant)**
  - **Smaller values in left subtree**
  - **Larger values in right subtree**

# Binary Search Tree – Insertion

## ■ Algorithm

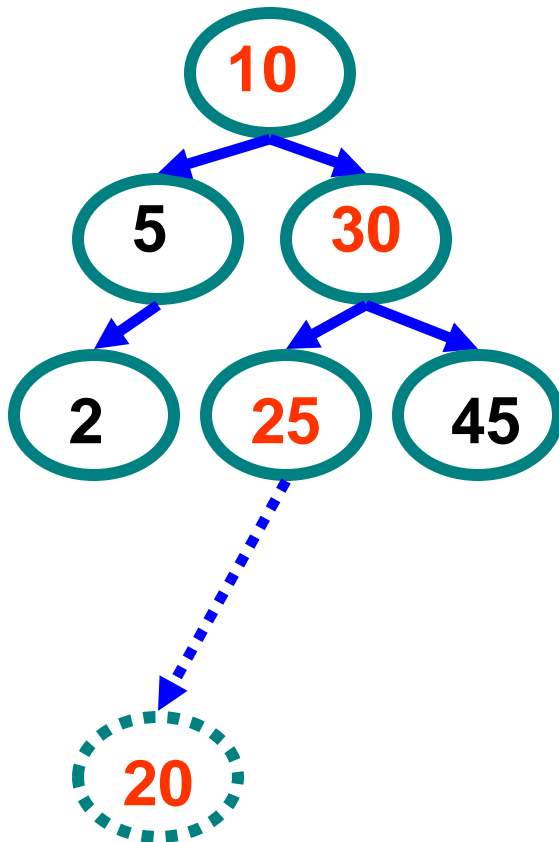
1. Perform search for value  $X$
2. Search will end at node  $Y$  (if  $X$  not in tree)
3. If  $X < Y$ , insert new leaf  $X$  as new left subtree for  $Y$
4. If  $X > Y$ , insert new leaf  $X$  as new right subtree for  $Y$

## ■ Observations

- $O(\log(n))$  operation for balanced tree
- Insertions may unbalance tree

# Example Insertion

■ Insert ( 20 )



$20 > 10$ , right

$20 < 30$ , left

$20 < 25$ , left

Insert 20 on left

# Binary Search Tree – Deletion

## ■ Algorithm

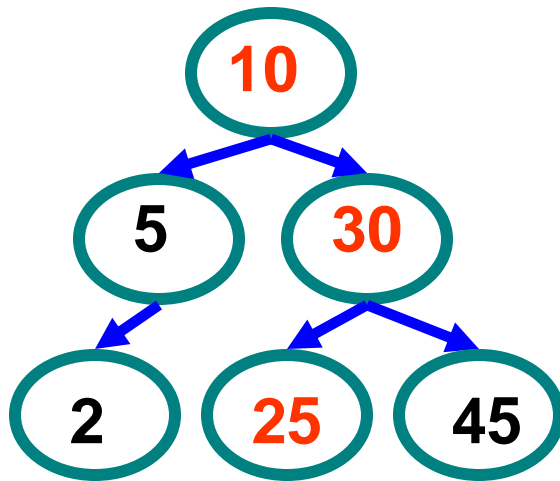
1. Perform search for value X
2. If X is a leaf, delete X
3. Else // must delete internal node
  - a) Replace with largest value Y on left subtree  
OR smallest value Z on right subtree
  - b) Delete replacement value (Y or Z) from subtree

## ■ Observation

- $O(\log(n))$  operation for balanced tree
- Deletions may unbalance tree

# Example Deletion (Leaf)

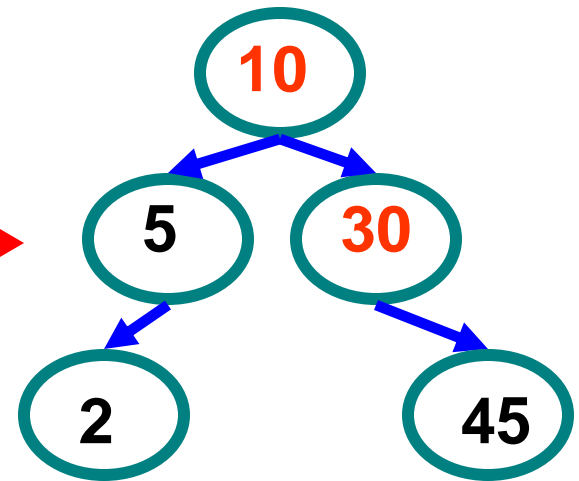
■ Delete ( 25 )



$25 > 10$ , right

$25 < 30$ , left

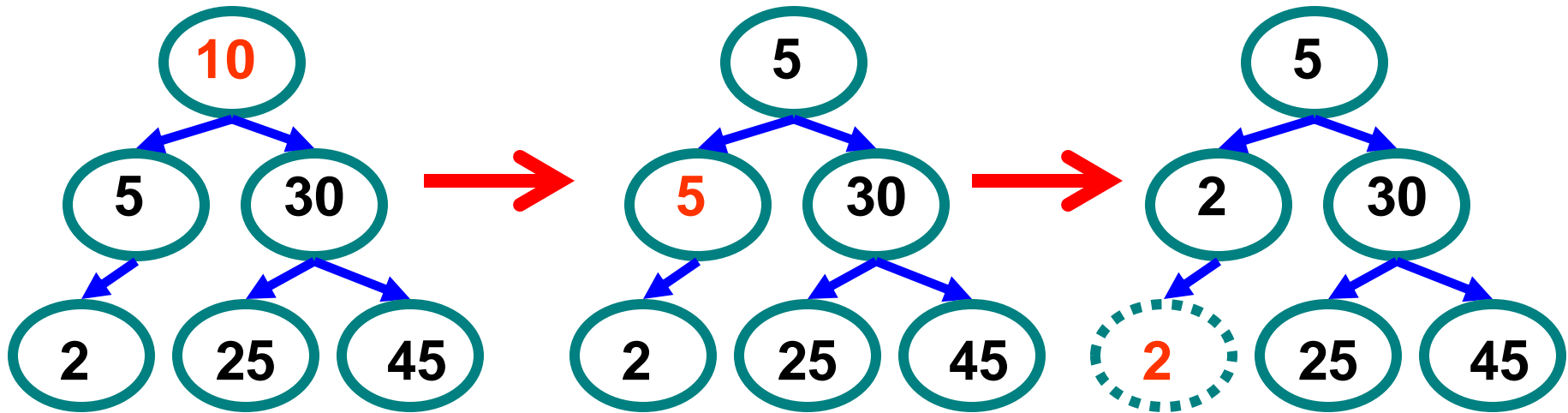
$25 = 25$ , delete





# Example Deletion (Internal Node)

■ Delete ( 10 )



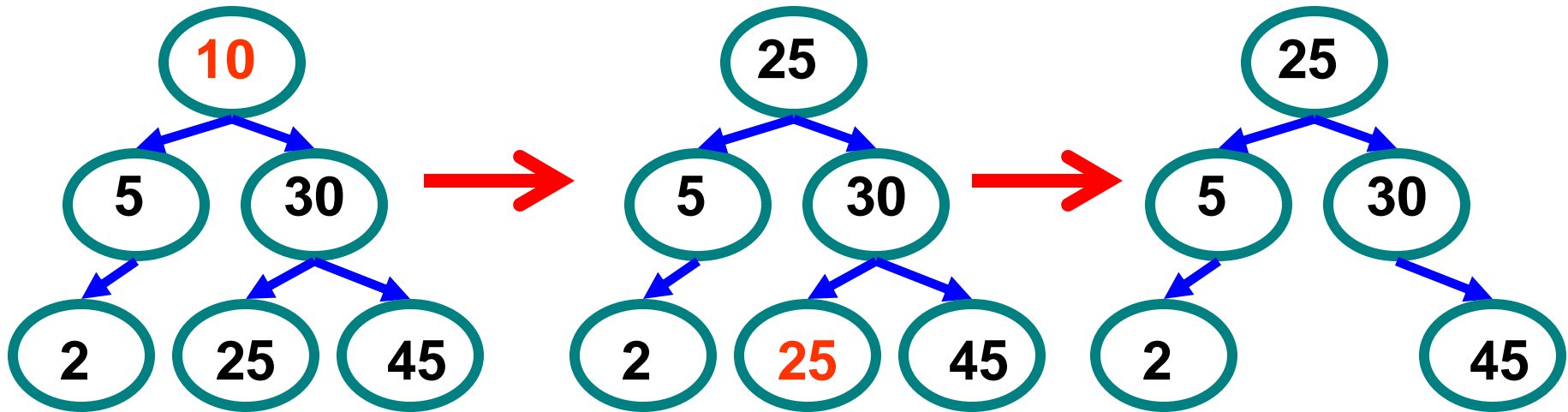
Replacing 10  
with **largest**  
value in left  
subtree

Replacing 5  
with **largest**  
value in left  
subtree

Deleting leaf

# Example Deletion (Internal Node)

■ Delete ( 10 )



Replacing 10  
with **smallest**  
value in right  
subtree

Deleting leaf

Resulting tree

# Building Maps w/ Search Trees

- Binary Search trees often used to implement **maps**
  - Each non-empty node contains
    - Key
    - Value
    - Left and right child
- Need to be able to compare keys
  - Generic type `<K extends Comparable<K>>`
    - Denotes any type K that can be compared to K's

# BST (Binary Search Tree) Implementation

- Implementing Tree using traditional approach
- Based on the BST definition below let's see how to implement typical BST Operations (constructor, add, print, find, isEmpty, isFull, size, height, etc.)

```
public class BinarySearchTree <K extends Comparable<K>, V> {  
    private class Node {  
        private K key;  
        private V data;  
        private Node left, right;  
        public Node(K key, V data) {  
            this.key = key;  
            this.data = data;  
        }  
    }  
    private Node root;  
}
```

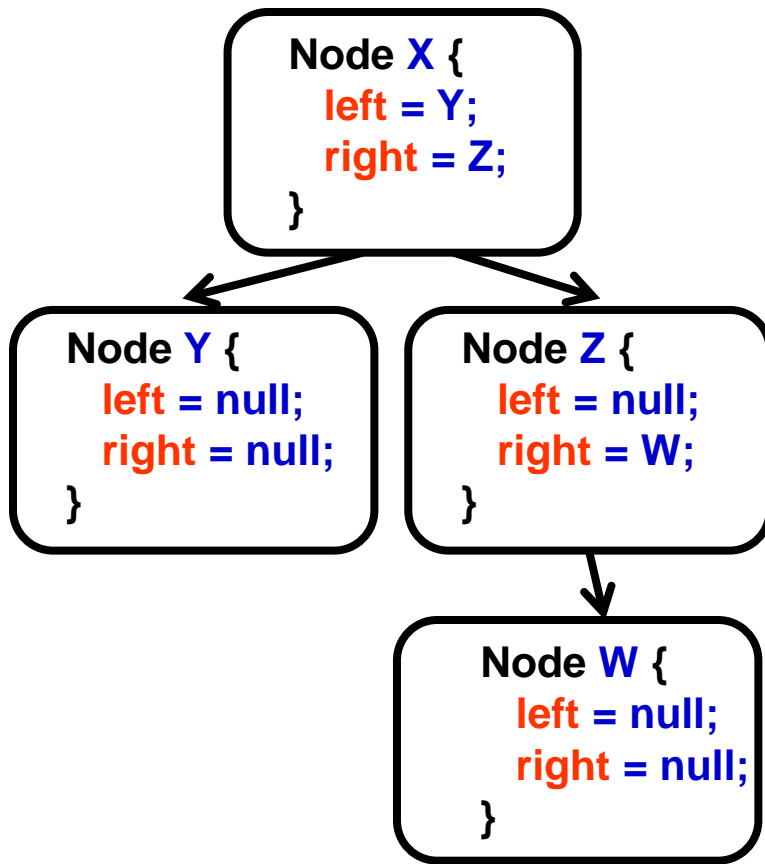
- See code distribution [BinaryTreeCode.zip](#)

# Polymorphic Binary Search Trees

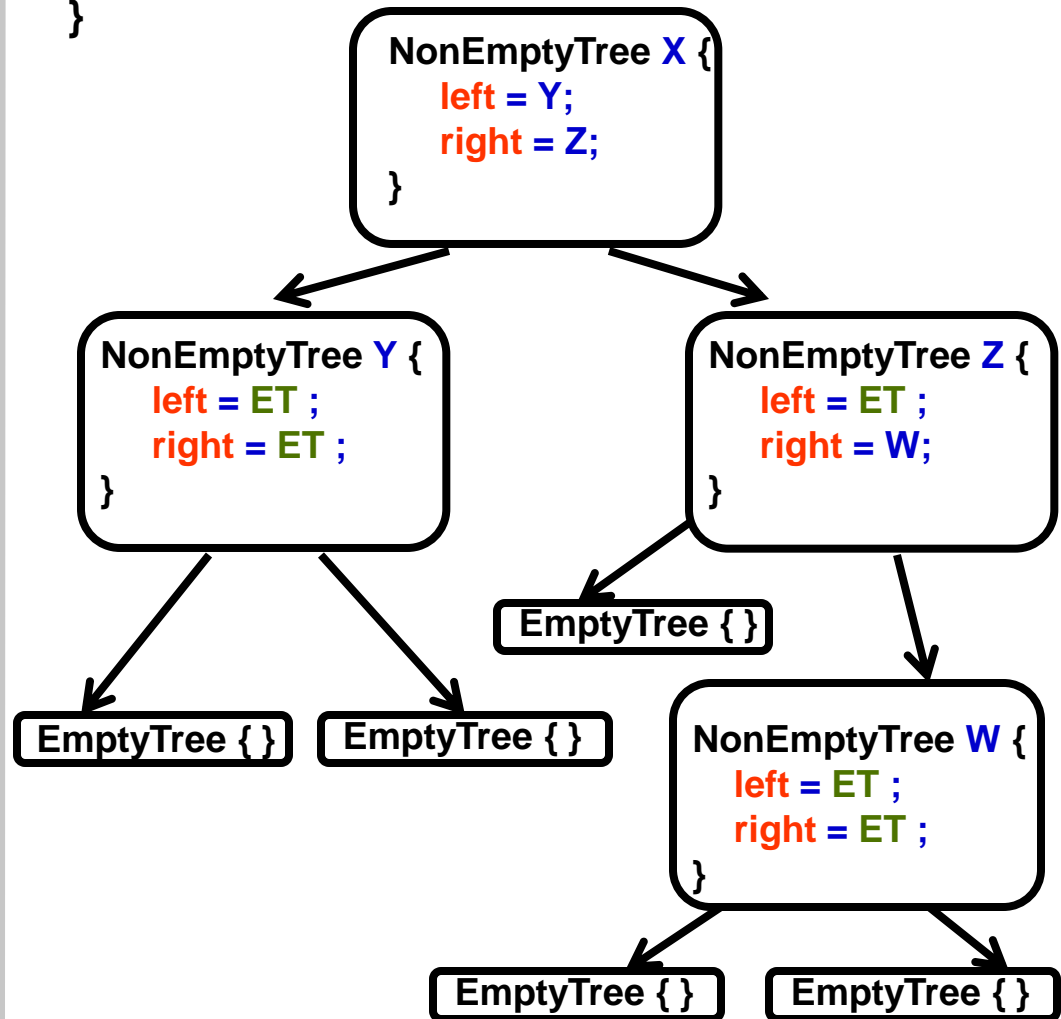
- **Second approach to implement BST**
- **What do we mean by polymorphic?**
- **Implement two subtypes of Tree**
  1. **EmptyTree**
  2. **NonEmptyTree**
- **Use EmptyTree to represent the empty tree**
  - **Rather than null**
- **Invoke methods on tree nodes**
  - **Without checking for null (IMPORTANT!)**

# Standard vs. Polymorphic Binary Tree

```
Class Node {  
    Node left, right;  
}
```



```
Class EmptyTree {}  
Class NonEmptyTree {  
    Tree left, right;  
}
```



# Polymorphic Binary Tree Implementation

Interface Tree {

    Tree insert ( Value data1 ) { ... }

}

Class EmptyTree implements Tree {

    Tree insert ( Value data1 ) { ... }

}

Class NonEmptyTree implements Tree {

    Value data;

    Tree left, right; // Either Empty or NonEmpty

    Tree insert ( Value data1 ) { ... }

}

# Singleton Design Pattern

## ■ Definition

- One instance of a class or value accessible globally

## ■ Where to use & benefits

- Ensure unique instance by defining class final
- Access to the instance only via methods provided

## ■ EmptyTree class will be a singleton class

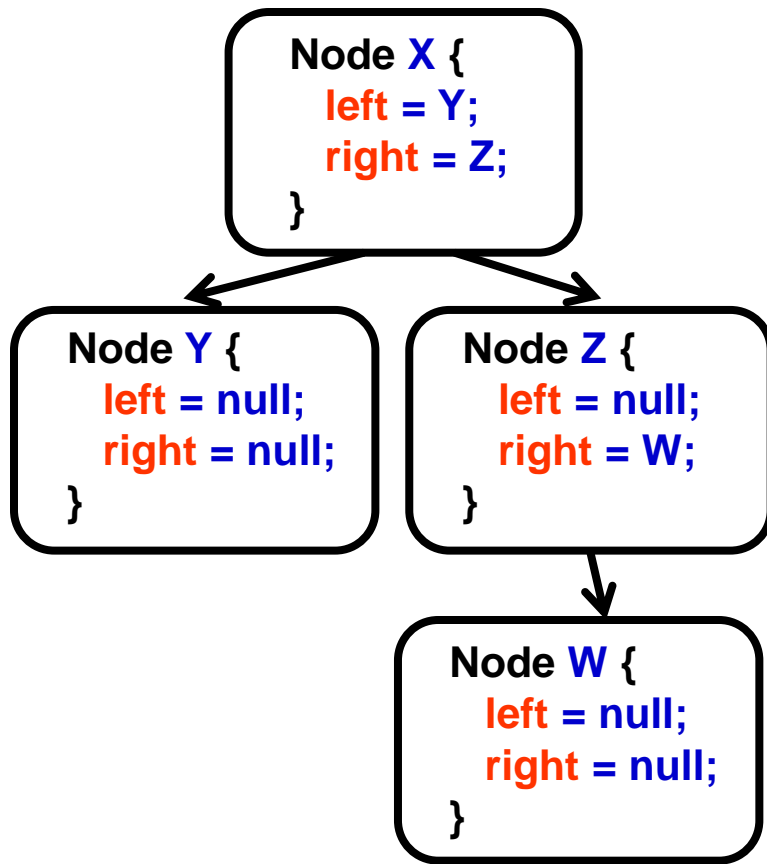


# Singleton Example

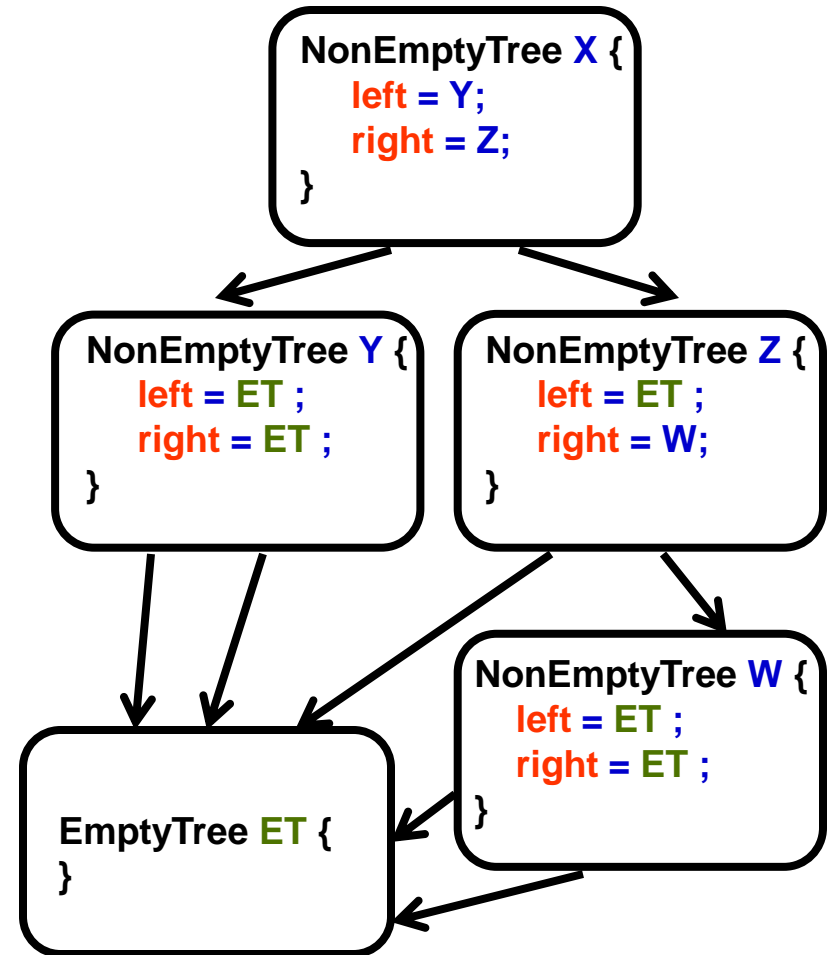
```
public final class MySingleton {  
    // declare the unique instance of the class  
    private static MySingleton uniq = new MySingleton();  
    // private constructor only accessed from this class  
    private MySingleton() { ... }  
    // return reference to unique instance of class  
    public static MySingleton getInstance() {  
        return uniq;  
    }  
}
```

# Using Singleton EmptyTree

```
Class Node {  
    Node left, right;  
}
```



```
Class EmptyTree {}  
Class NonEmptyTree {  
    Tree left, right;  
}
```



# Polymorphic List Implementation

- Let's see a polymorphic list implementation
- See code distribution `PolymorphicListCode.zip`